



Quantum Cryptography

Marshall Roth
March 9, 2007



Overview

- Current Cryptography Methods
- Quantum Solutions
- Quantum Cryptography
- Commercial Implementation

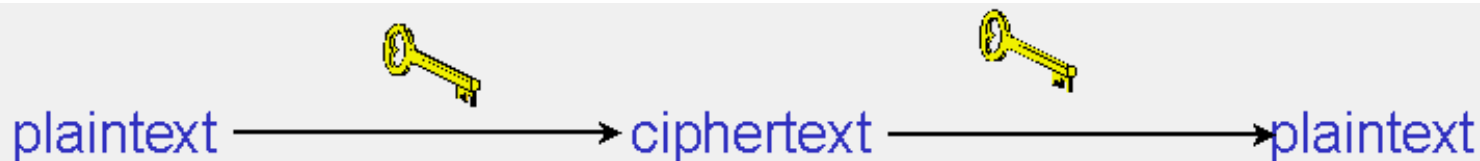


Cryptography algorithms:

Symmetric – encrypting and decrypting key are identical (Data Encryption Standard, Rivest Ciphers)

Asymmetric – encrypting and decrypting keys differ (Elliptical Curve; Rivest, Shamir, Adleman)

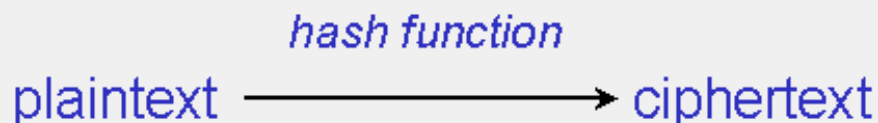
Hash – no decryption by design, meant to uniquely identify a message such as a password (Message Digest)



A) Secret key (symmetric) cryptography. SKC uses a single key for both encryption and decryption.



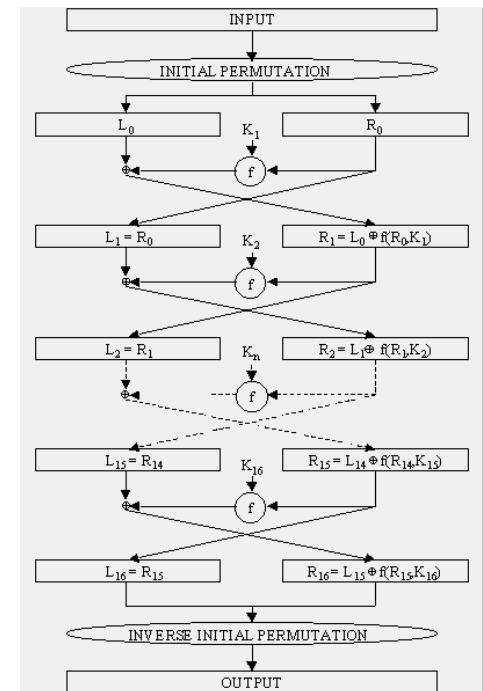
B) Public key (asymmetric) cryptography. PKC uses two keys, one for encryption and the other for decryption.



C) Hash function (one-way cryptography). Hash functions have no key since the plaintext is not recoverable from the ciphertext.

Symmetric Key Distribution

- RC5 and others takes sufficiently long decrypt (72 bits with distributed computing ~1000 years for RC5)
- How do we securely distribute keys?
- Some methods work on simple binary addition:
$$s = m \oplus k, m = s \oplus k = m \oplus k \oplus k$$
- Others, such as DES, shuffle blocks of information





Asymmetric Key Distribution

- Rivest, Shamir, Adelman (RSA) use the property of factoring a large number in terms of primes is sufficiently complex with classical computers.
- Elliptical Curves make use of another sufficiently complex classical problem of calculating the discrete logarithm.
- Codes can be broken more readily than symmetric keys (72 bits sym ~ 2048 bits asym)



RSA Algorithm

- Pick two large prime numbers p and q and calculate the product $N = pq$, $\phi = (p - 1)(q - 1)$
- Choose a number that is co-prime with ϕ , c
- Find a number d to satisfy $cd = 1 \pmod{\phi}$, using a method such as Euclid's algorithm
- Using your plaintext, a , the ciphertext is encoded as $b = a^c \pmod{N}$
- To retrieve the plaintext, $a = b^d \pmod{N}$
- The numbers N and c are made public, so anyone can encrypt information, but only someone with d can retrieve the plaintext

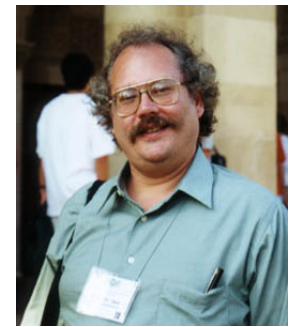


Example

- Plaintext $a = 123$
- $p = 61$ and $q = 53$
- $N = pq = 3233$
- $\phi = (p - 1)(q - 1) = 3120$
- Pick a coprime of ϕ , $c=17$
- Find d such that $cd = 1 \pmod{\phi}$, $d=2753$
- Encode with $a^c \pmod{N}$, in this case $123^{17} \pmod{3233} = 855$
- Decode message by evaluating $b^d \pmod{N}$, in this case $855^{2753} \pmod{3233} = 123$



Enter Shor's Algorithm



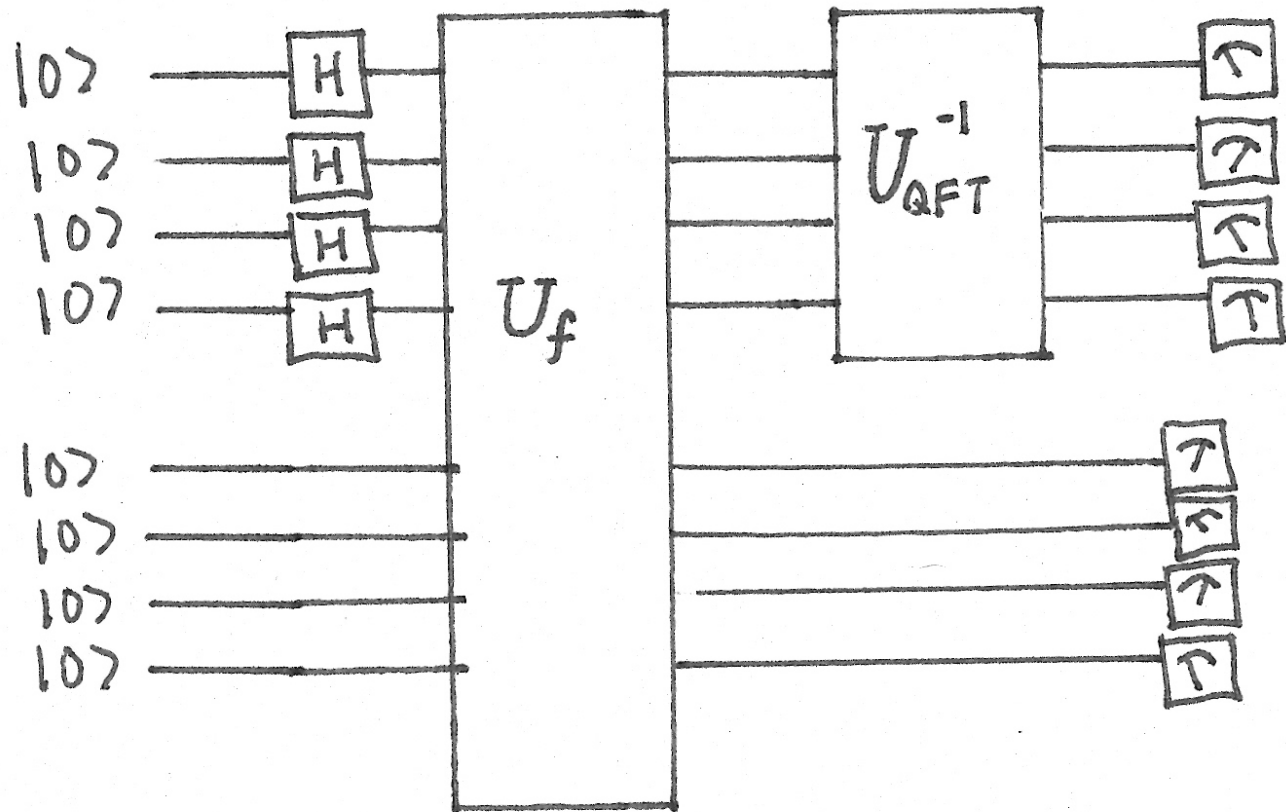
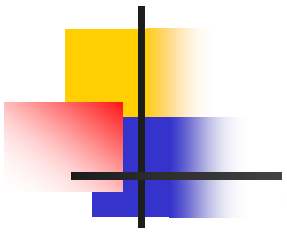
- Let $f(x) = b^x \bmod N$, if we can find some r that $f(x) = f(x+r)$, then we can find a number d' such that $cd' = 1 \bmod r$
- The value d' works like the decoding value we calculated from $cd = 1 \bmod \phi$
- In addition, using different values for $b < N$, we can determine the prime components of N

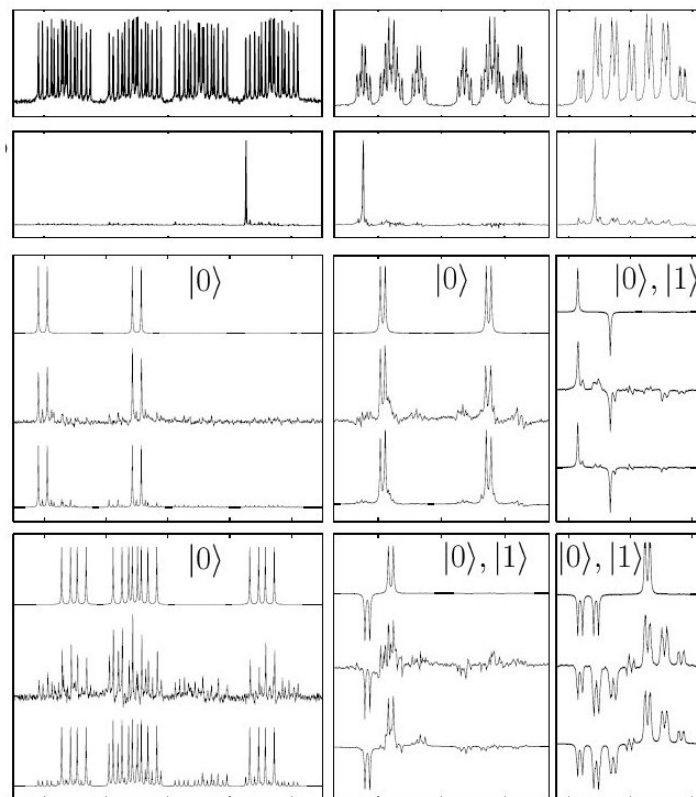
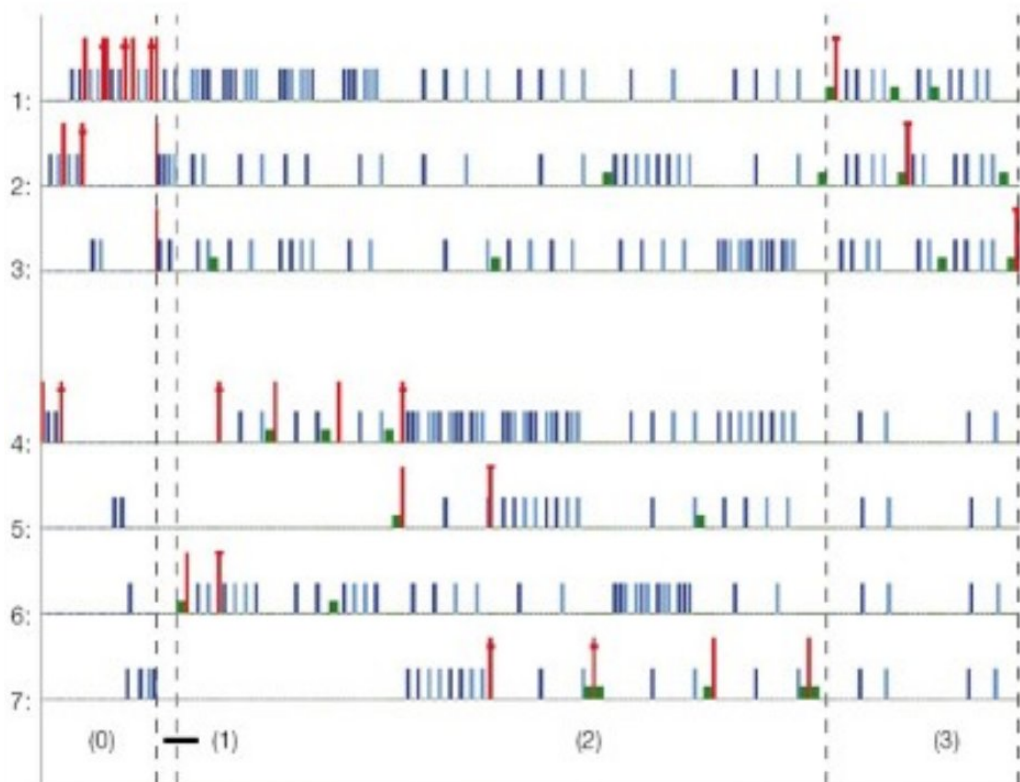
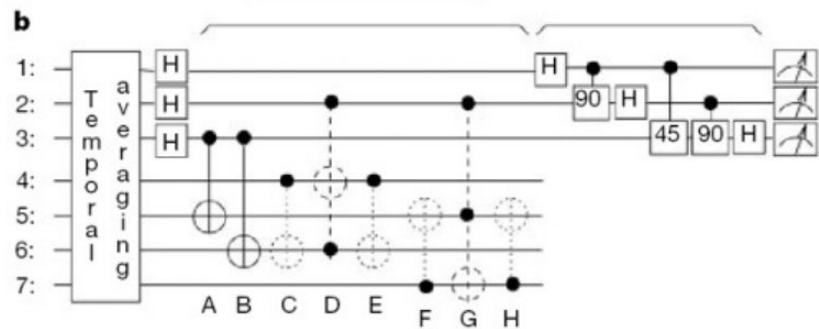
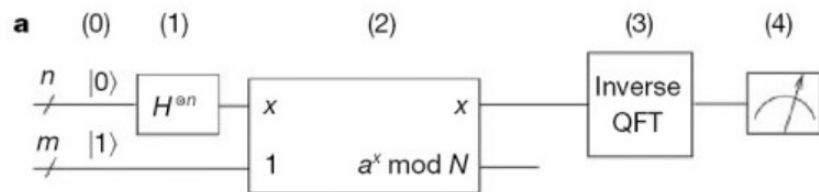
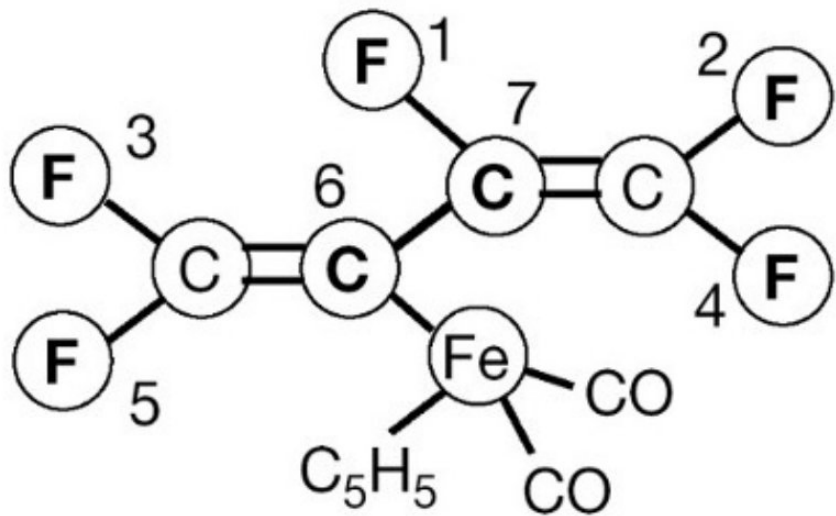
What's the quantum algorithm?

- Initialize $\log_2 N$ qubits an equal superposition state (input qubits)

$$|\psi\rangle = \frac{1}{2^{n/2}} (|000\dots000\rangle + |000\dots001\rangle + \dots + |111\dots111\rangle)$$

- Using $\log_2 N$ more qubits, enact $f(\psi)$ on them while retaining the state ψ in the input state (output qubits)
- Apply the quantum fourier transform on the ψ portion of the circuit
- Measure the input and output qubits $(y, f(x_0))$, with high probability you will measure an of $f(x_0 + (y/N)r)$ where y/N is close to an integer





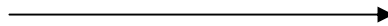


So we've broken it, now what?

- In general symmetric keys are harder to crack but tough to distribute, while asymmetric keys are easy to distribute but easier to crack
- Start thinking about using quantum systems to implement cryptography
- Restrictions on polarization bases measurements ($\uparrow \rightarrow \searrow \swarrow \sigma^+ \sigma^-$)
- Restrictions on state duplication
- Very easy to create state perturbation

BB84 (Bennett and Brassard)

- We have two parties, Alice and Bob who want to securely distribute their symmetric key over a public channel





BB84



- Alice randomly chooses one of two orientations from two bases to measure in: (for spin $\frac{1}{2}$ situation analogous to z-basis, and x-basis)
- Alice then assigns the value of 0 and 1 in each basis (up-z and up-x = 0, down-z and down-x = 1)
- Alice sends a state from one of the four bases at random, and Bob selects (with his own random generator) a basis (x or z) to measure in
- If they choose the same basis, they will agree with 100% probability, if they choose a different basis they will have no way of correlating the results (error rate $\sim 25\%$)



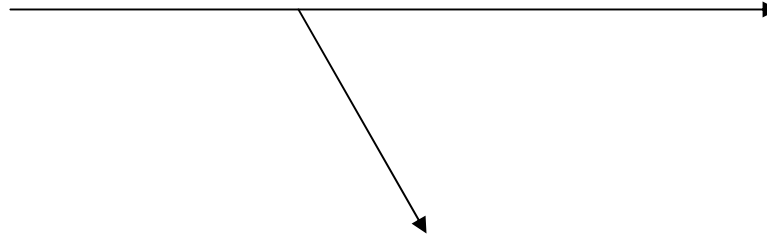


BB84

- In order to verify the transmitted information, Alice and Bob decide which bits can be kept and which bits will need to be retransmitted
- The correlated measurements will only be in compatible bases (both x or z)

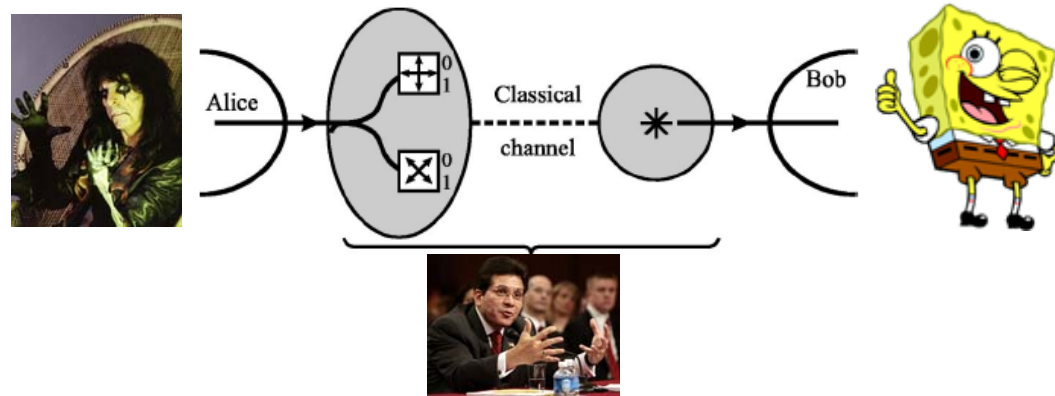
<http://monet.mercersburg.edu/henle/bb84/demo.php>

Eavesdropping



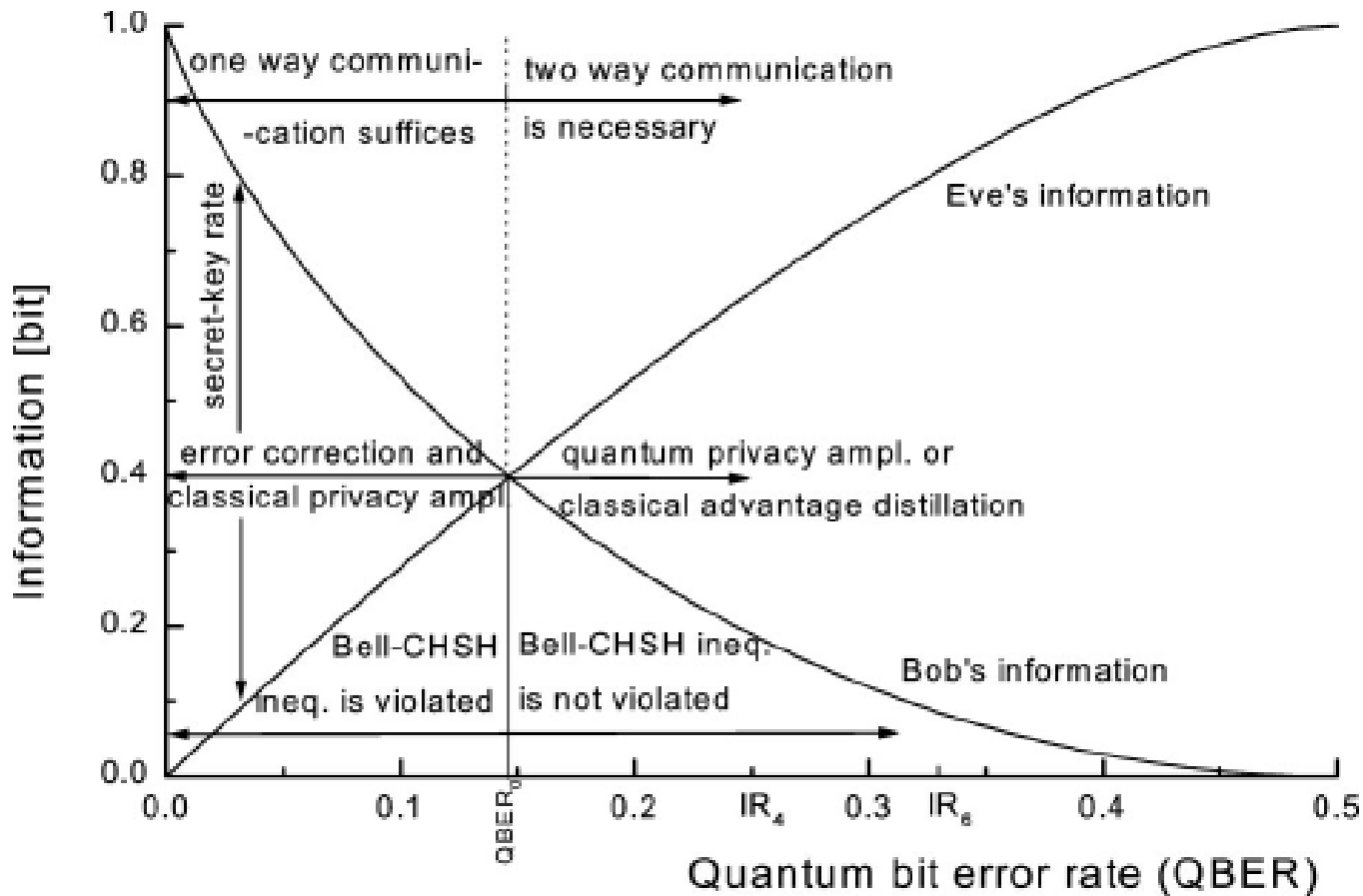
Multiple Photon Attack

- Eve can attack an optical channel by measuring multiple photon signals with a PBS and recreating the signals



2/3 of the time Eve can recreate the original state and send it to Bob, the rest of the time she introduces an error rate of 1/6

Eavesdropping Thresholds





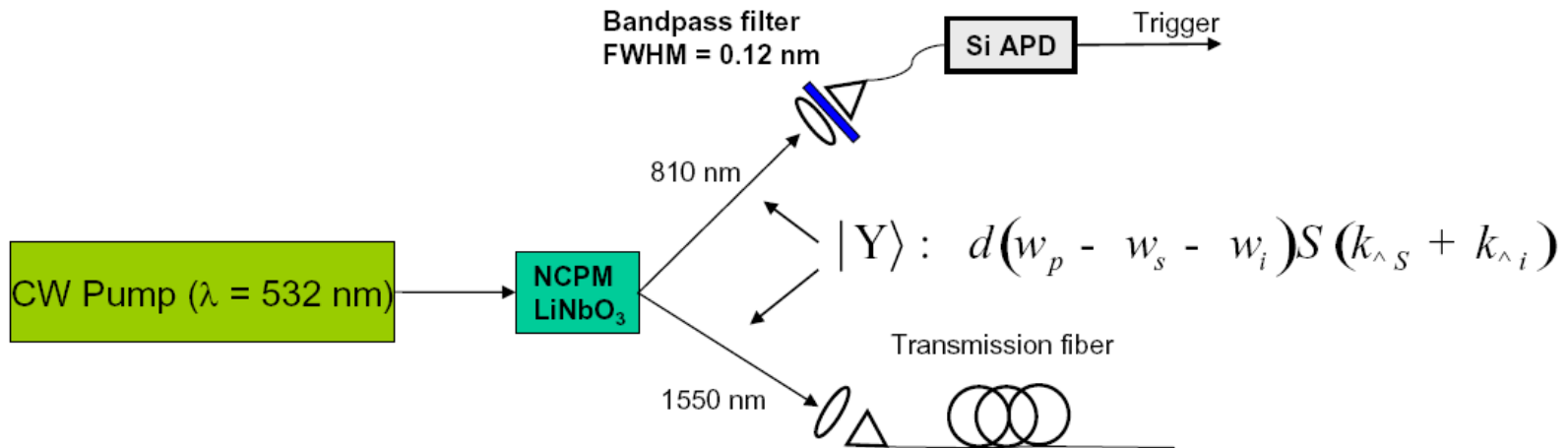
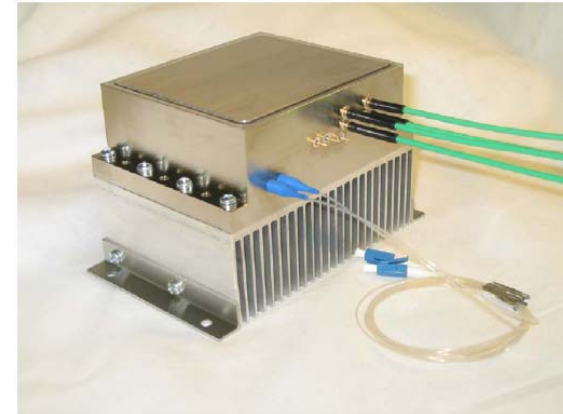
Commerical Quantum Crypto Systems

- Magiq: Currently have an implementation of a secure quantum network, the QPN 7505
- Works on a single photon source, and can transmit up to about 75 km with reasonable loss

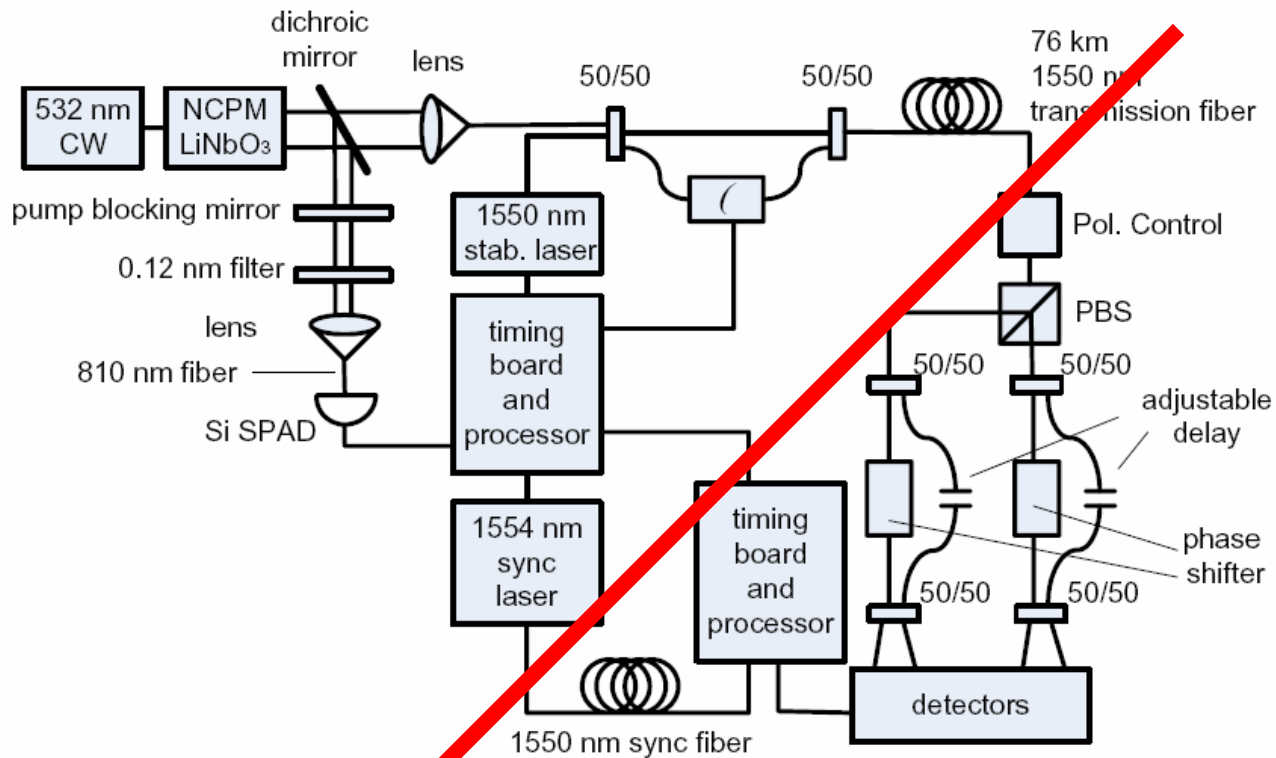
Price? \$97000

Single Photon Manipulation

Entanglement occurs in the time-frequency domain, there is a high probability that a single photon is produced, and low probability of multiple photons

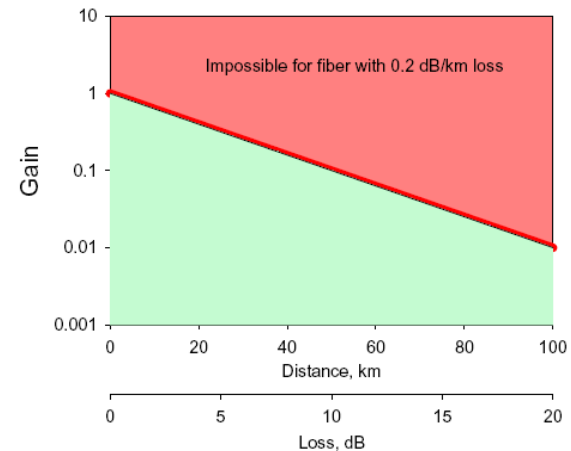
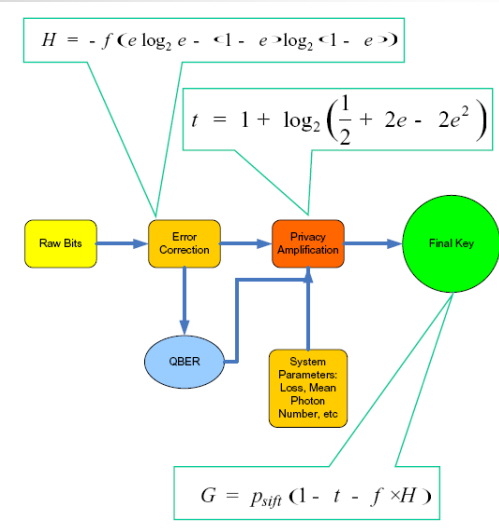


Single Photon setup

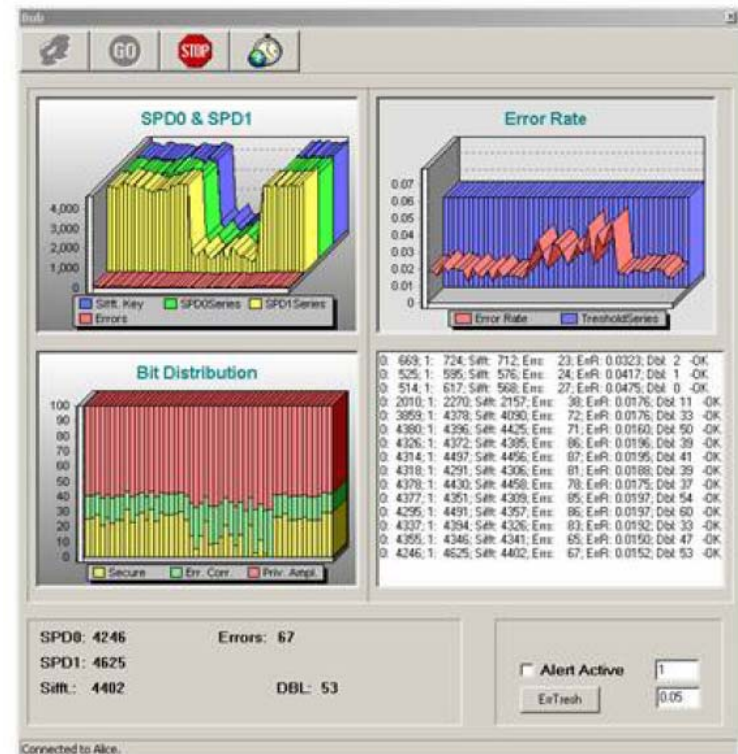
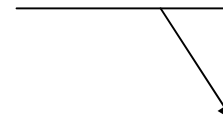


Fidelity Considerations

- The system tries to maximize G , the probability of transmitting a secure bit with a single initial pulse
- This attenuates about 10dB for every 50 km of transmission

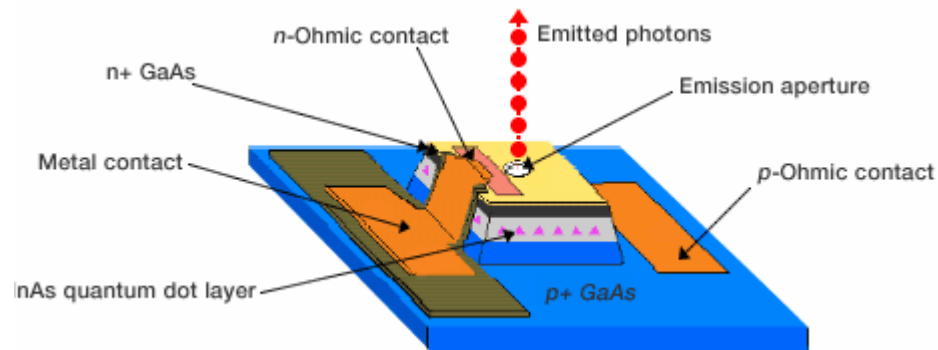


Eavesdropping!



Other systems

- Id quantique:
 - Vectis – Their crypto system, uses typical QKD and AES (Advanced Encryption Standard)
 - Quantis – Random number generator, based on standard 50/50 polarization probabilities (4 Mbit/s number generation) PCI hardware
- Toshiba Research – Cambridge, QKD and single photon emission with quantum dots





Sources

- "*Quantum Cryptography*"; Gisin, Ribordy, Tittel, Zbinden; 2002
- "*Secure Communication with single photons*"; A. Trifonov; 2005
- "*An Overview of Cryptography*"; Gary Kessler; 2007
- "*Applications of Quantum Cryptography in Government Classified, Business, and Financial Communications*"; Audrius Berzanskis; 2005
- "*Quantum key distribution over 122 km of standard telecom fiber*"; Gobby, Yuan, Shields; 2004