

a)

**TmemoryManager's Constructor**

```
FreeLinksList <- get TmyList object reference (call to TmyList())
pointer index <- 0
while( index LESS THAN MaxListSize-SizeOfLink )
  call TmyList's insertElement(-1, index) for FreeLinkList's
  index <- index + SizeOfLink
```

**Running Time Calculation**

$$\frac{(\text{MaxListSize}-\text{SizeOfLink})}{\text{SizeOfLink}} = \text{MaxListSize}/\text{SizeOfLink}$$

index = 0

**TmyList's Constructor**

```
head <- tail <- null
```

**TmyList's public void insertElement( int key, int freeLinkAddress )**

```
// If this is really a free link, set the key. If this is an empty
// list set the head to point to the new link, else add the item to
// the end of the existing list. After adding the item set the tail
// pointer and set the new item's next pointer to null
if( freeLinkAddress NOT EQUAL null )
  dereference freeLinkAddress <- key
  if( listIsEmpty()
      head <- freeLinkAddress
  else
      dereference tail's next <- freeLinkAddress
  tail <- freeLinkAddress
  dereference tail's next <- null
```

**Worst and Average Case**

Cost	Times
1	1
1	1
1	MaxListSize/SizeOfLink (see below)
1	MaxListSize/SizeOfLink
1	MaxListSize/SizeOfLink

**Worst and Average Case**

Cost	Times
1	1

**Worst Case**

Cost	Time
1	1
1	1
1	1
1	1/2
1	1/2
1	1
1	1

To determine the worst case running time of TmemoryManager(), we must take into account the running time of TmyList() and insertElement(...).

TmyList()'s  $T(n) = 1 * 1 = 1$

insertElement ()'s  $T(n) = (1 * 1) + (1 * 1) + (1 * 1) + (1 * 1/2) + (1 * 1/2) + (1 * 1) + (1 * 1) = 1 + 1 + 1 + 1/2 + 1/2 + 1 + 1 = 6$

TmemoryManager's  $T(n) = (\text{TmyList}'\text{s } T(n) * (1 * 1)) + (1 * 1) + (1 * \text{MaxListSize}/\text{SizeOfLink}) + (\text{insertElement } ()'\text{s } T(n) * (1 * \text{MaxListSize}/\text{SizeOfLink})) + 1 * \text{MaxListSize}/\text{SizeOfLink}$

TmemoryManager's  $T(n) = (1 * 1) + 1 + \text{MaxListSize}/\text{SizeOfLink} + (6 * \text{MaxListSize}/\text{SizeOfLink}) + \text{MaxListSize}/\text{SizeOfLink}$

TmemoryManager's  $T(n) = 2 * \text{MaxListSize} / \text{SizeOfLink} + 2$   
 $O(\text{MaxListSize} / \text{SizeOfLink})$

The average and worst case run time is the same, because you are guaranteed to execute the statements based on MaxListSize and the SizeOfLink.

In our case MaxListSize = 30

```
index   : 0  1  2  3  4  5  6  7  8  9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25 26 27 28 29
value   : -1  2 -1  4 -1  6 -1  8 -1 10 -1 12 -1 14 -1 16 -1 18 -1 20 -1 22 -1 24 -1 26 -1 28 -1 -1
key/next: k  n  k  n  k  n  k  n  k  n  k  n  k  n  k  n  k  n  k  n  k  n  k  n  k  n  k  n  k  n
```

b)

**TmyList's int removeLinkFromList( int key )**

```
pointer previous <- null
pointer current  <- head

// if match is first element, reset fHead, remove and return link
if(( NOT listIsEmpty()) AND ( key EQUALS dereference of current ))
    call removeFirstLinkFromList()
    return current

// not found so, walk through the rest of the list
previous <- current
current  <- dereference current's next
while( current NOT EQUAL null )
    // If a match is found, set the previous pointer's next pointer
    // to the value of current's next pointer, therefore unlinking
    // current.  If current was the last link then reset the tail
    // pointer and return the unlinked value
    if( key EQUALS dereference of current)
        dereference previous' next <- dereference current's next
        if( dereference previous' next EQUALS null )
            tail <- null
            return current
        previous <- current
        current  <- dereference current's next

// if key is not found, an invalid address is returned.
// and the content of the list is not altered
return current
```

**TmyList's public void insertElement( int key, int freeLinkAddress )**

```
// If this is really a free link, set the key.  If this is an empty
// list set the head to point to the new link, else add the item to
```

```

// the end of the existing list. After adding the item set the tail
// pointer and set the new item's next pointer to null
if( freeLinkAddress NOT EQUAL null )
    dereference freeLinkAddress <- key
    if( listIsEmpty()
        head <- freeLinkAddress
    else
        dereference tail's next <- freeLinkAddress
    tail <- freeLinkAddress
    dereference tail's next <- null

```

c)

### TmemoryManager's public int nextFreeLink()

```

// returns the address of the next free linkEntry
// if there are no more free ones, return kInvalidAddress.
return result of calling TmyList's removeFirstLinkFromList() for FreeLinksList

```

<u>Worst Case</u>	
Cost	Time
1	1

### TmyList's public int removeFirstLinkFromList()

```

// delete the first link from the list, if list is empty
// an invalid address is returned.
pointer current <- head
if( NOT listIsEmpty()
    head <- dereference current's next
    // if first and only element then reset tail
    if( listIsEmpty()
        tail <- null
return current

```

<u>Worst Case</u>	
Cost	Time
1	1
1	1
1	1
1	1
1	1
1	1

### Running Time Calculation

To determine the worst case running time of nextFreeLink(), we must take into account the running time of removeFirstLinkFromList().

removeFirstLinkFromList()'s  $T(n) = (1*1) + (1*1) + (1*1) + (1*1) + (1*1) + (1*1) = 1+1+1+1+1+1=6$

nextFreeLink()'s  $T(n) = \text{removeFirstLinkFromList()'s } T(n) * (1*1) + = 1 * 6 = 6$

$O(1)$

### TmemoryManager's public void returnFreeLink(int linkAddress)

```

// linkAddress is a linkEntry that is not being used anymore.
// add the free element to the free Linked List
call TmyList's insertElement(-1, linkAddress) for FreeLinkList

```

<u>Worst Case</u>	
Cost	Time
1	1

### TmyList's public void insertElement( int key, int freeLinkAddress )

See part a) above for analysis of insertElement(...)

### Running Time Calculation

To determine the worst case running time of `returnFreeLink()`, we must take into account the running time of `insertElement()`.  
`insertElement()`'s  $T(n) = (1*1) + (1*1) + (1*1) + (1*1/2) + (1*1/2) + (1*1) + (1*1) = 1+1+1+1/2+1/2+1+1 = 6$   
`returnFreeLink()`'s  $T(n) = \text{insertElement}()$ 's  $T(n) * (1*1) = 6 * 1 = 6$   
 $O(1)$

d)  
**TmyQueue's Constructor**

`fQueue <- get TmyList object reference (call to TmyList())`

<u>Worst Case</u>	
Cost	Time
1	1

**See part a) for TmyList() Analysis**

$T(n) = \text{TmyList}()$ 's  $T(n) * (1*1) = 1 * 1 = 1$   
 $O(1)$

**TmyQueue's public void enqueue( int key, int freeLinkAddress)**

`call insertElement(key, freeLinkAddress) for the queue`

<u>Worst Case</u>	
Cost	Time
1	1

**See part a) above for analysis of insertElement(...)**

**Running Time Calculation**

To determine the worst case running time of `enqueue()`, we must take into account the running time of `insertElement()`.

`insertElement()`'s  $T(n) = (1*1) + (1*1) + (1*1) + (1*1/2) + (1*1/2) + (1*1) + (1*1) = 1+1+1+1/2+1/2+1+1 = 6$

`enqueue()`'s  $T(n) = \text{insertElement}()$ 's  $T(n) * (1*1) = 6 * 1 = 6$

$O(1)$

**TmyQueue's public int dequeue()**

`return the result of calling TmyList's removeFirstLinkFromList()`

<u>Worst Case</u>	
Cost	Time
1	1

**See part c) for analysis of removeFirstLinkFromList()**

To determine the worst case running time of `dequeue()`, we must take into account the running time of `removeFirstLinkFromList()`.

`removeFirstLinkFromList()`'s  $T(n) = (1*1) + (1*1) + (1*1) + (1*1) + (1*1) + (1*1) = 1+1+1+1+1+1=6$

`dequeue()`'s  $T(n) = \text{removeFirstLinkFromList}()$ 's  $T(n) * (1*1) = 6 * 1 = 6$

$O(1)$