LKB Formalism Lab 1 questions

Ling 567 January 6, 2010

Overview

- Type hierarchies, inheritance, unification
- Typed feature structures, subsumption, unification
- Type constraints, making typed feature structures wellformed
- Notational conventions
- Grammar rules in the LKB
- Lab 1 questions

tdl and typed feature structures

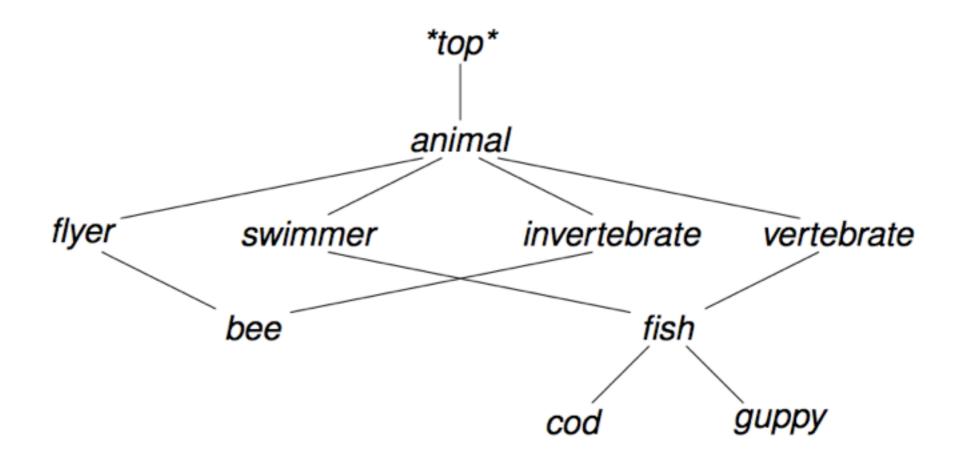
- •tdl = type description language
- •.tdl files encode type descriptions.
- The LKB reads in the tdl files and compiles the type descriptions into a well-formed type hierarchy.
- •NB: Actual trees are not subject to the constraint that they be fully specified, but they must be well-typed (all features appropriate for a type are present, though types need not be maximally specific).

Properties of our type hierarchies

- Unique top: All types ultimately inherit from one top node
- No cycles: No path through the hierarchy from a type to itself
- Unique greatest lower bounds (glbs): Any two types in the hierarchy are either incompatible (share no descendants) or have a unique most general subtype
- Closed world: All types that exist have a known position in the hierarchy
- Compatibility: Two compatible types unify to their glb

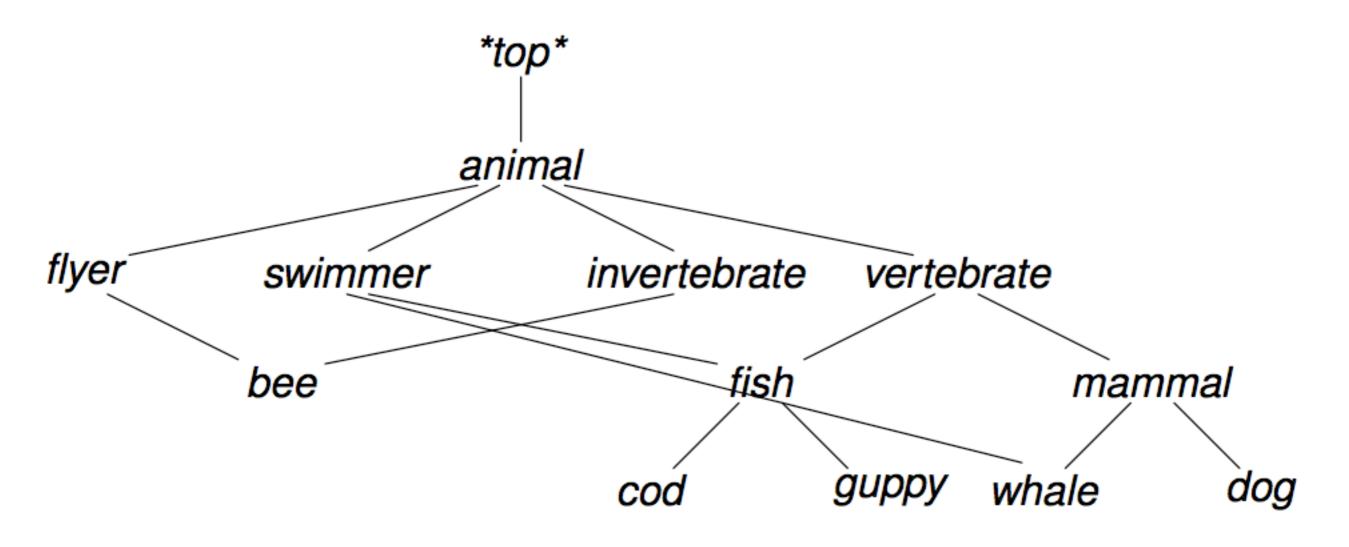
Multiple inheritance and unification

- flyer and swimmer are incompatible (no common descedants)
- flyer and bee unify to subtype (hierarchical relationship)
- flyer and invertebrate unify to glb (bee)



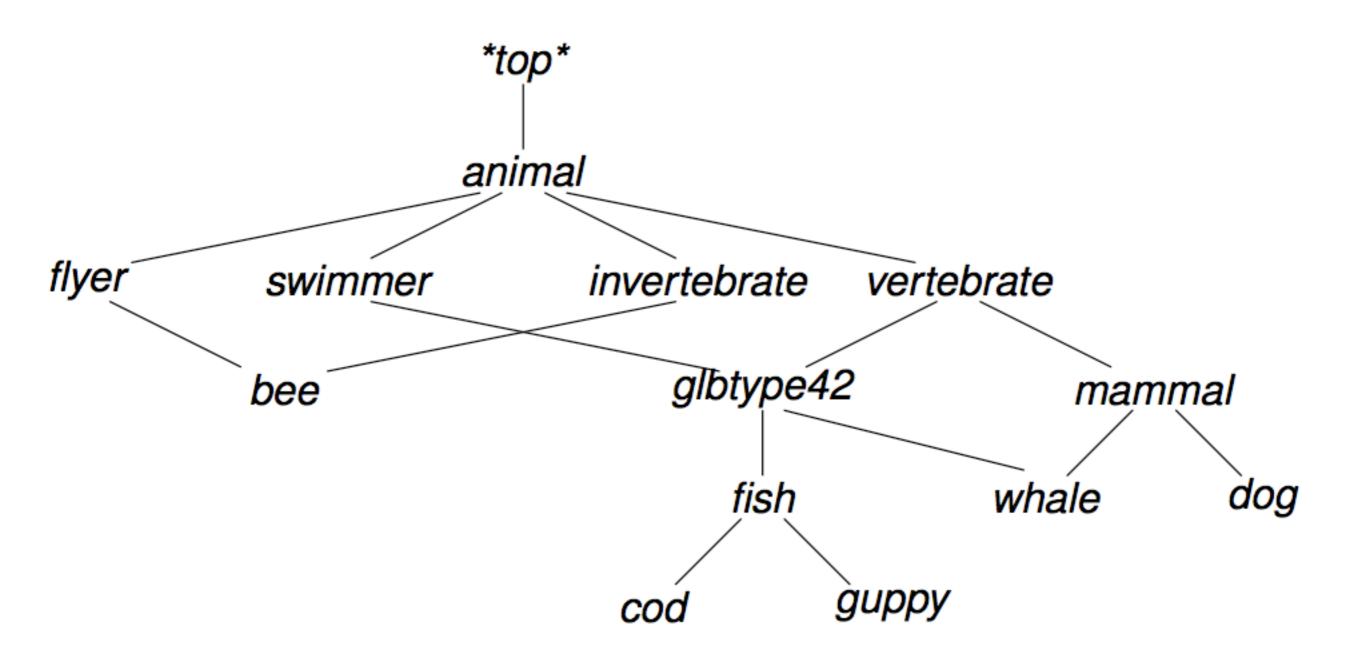
An invalid type hierarchy

- *swimmer* and *invertebrate* have two common subtypes: fish and whale
- fish and whale are incomparable in the hierarchy: glb condition is violated



Fixing the type hierachy

The LKB introduces glb types as required



Properties of typed feature structures

- Finiteness: A typed feature structure has a finite number of nodes
- Unique root and connectedness: A tfs has a unique root parent; all other nodes have at least one parent
- No cycles: No node has an arc that points back to the root node or to another node that intervenes between the node itself and the root
- Unique features: Any node can any (finite) number of outgoing arcs, but the arc labels (i.e., features) must be unique within each node
- Typing: Each node has a single type which is defined in the hierarchy

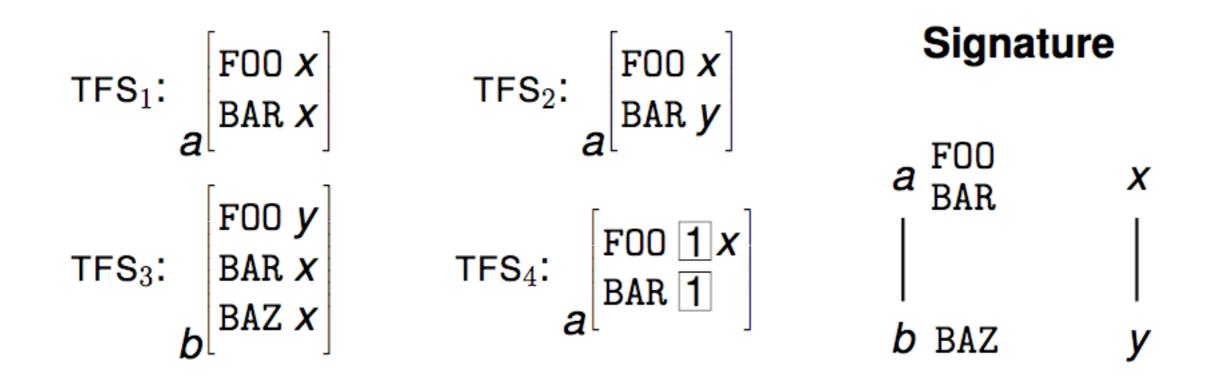
tdl example

type := supertype1 & supertype2 & [FEAT1 val1, FEAT2 val2 & [FEAT3 #same, FEAT4 #same]].

Typed feature structure subsumption

- •tfss can be partially ordered by information content
- a more general structure is said to subsume a more specific one
- *top* is the most general feature structure, while ⊥ is inconsistent
- Feature structure *F* subsumes feature structure *G* iff: (1) if path *p* is defined in *F* then *p* is also defined in *G* and the type of the value of *p* in *F* is as supertype or equal to the value of *p* in *G*, and (2) all paths that are reentrant in *F* are also reentrant in *G*.

Subsumption examples

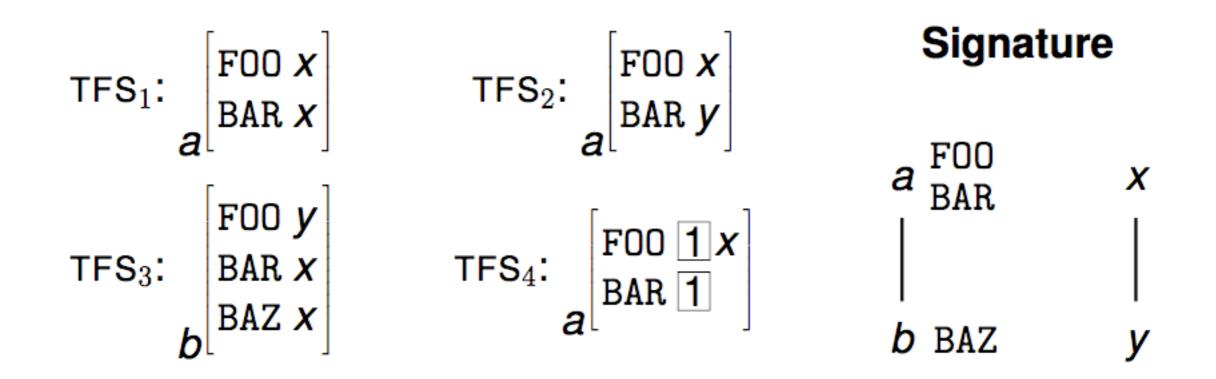


Which tfss subsume which other tfss?

Typed Feature Structure Unification

- Decide whether the two typed feature structures are compatible
- Determine the combination of the two tfss which gives the most general feature structure which retains all of the information they each individually contain
- Unification monotonically combines information from both 'input' tfss
- The unification of *F* and *G* is the most general tfs that is subsumed by both *F* and *G* (if it exists).

Unification examples



What is the unification of TFS1&2? 1&3? 3&4?

Type constraints and appropriate features

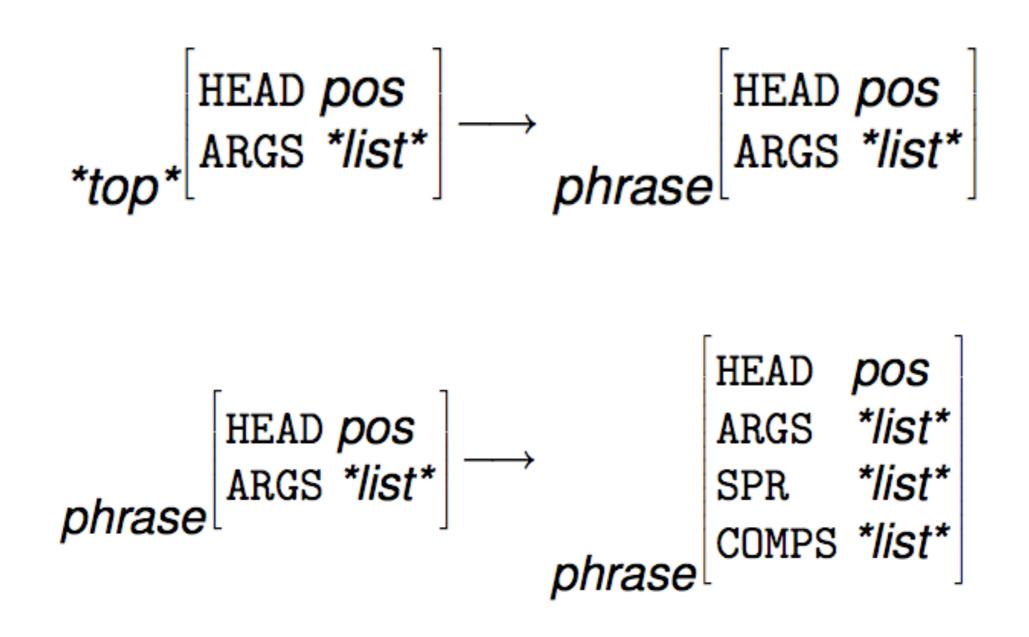
- Well-formed tfss satisfy all type constraints from the type hierarchy
- Type constraints are typed feature structures associated with a type
- The top-level features of a type constraint are its appropriate features

type	constraint	appropriate features
ne-list	FIRST *top* REST *list*	FIRST and REST

Type inference: Making a tfs well-formed

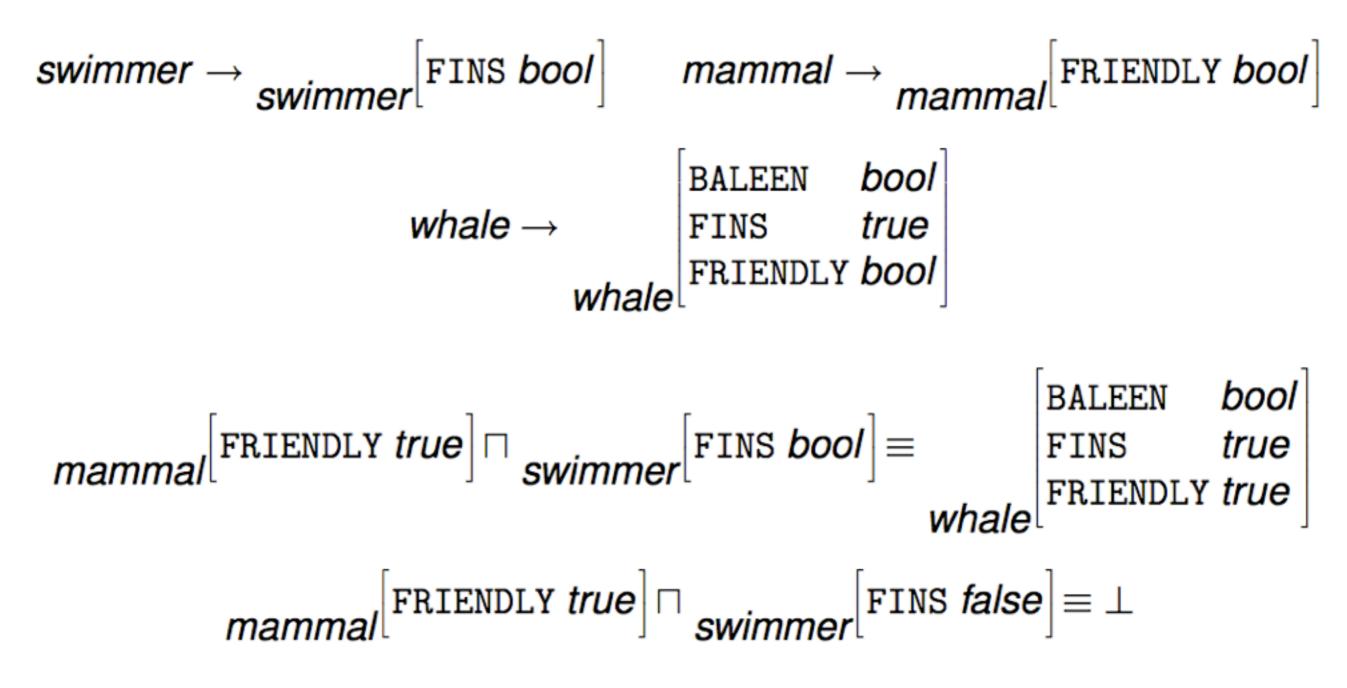
- Apply all type constraints to convert tfs to well-formed tfs
- Determine most general well-formed tfs subsumed by input tfs
- Specialize all types so that all features are appropriate
- Expand all nodes with the type constraint of the type on that node

Examples



More interesting well-formed unification

Type Constraints Associated to Earlier animal **Hierarchy**



Recursion in the type hierachy

 Type hierarchy must be finite after type inference; illegal type constraint:

list := *top* & [FIRST *top*, REST *list*].

Needs additional provision for empty lists; indirect recursion:

```
*list* := *top*.
*ne-list* := *list* & [ FIRST *top*, REST *list* ].
*null* := *list.
•Recursive types allow for parameterized list types:
*s-list* := *top*.
*s-ne-list* := *ne-list* & *s-list* &
[ FIRST *top*, REST *list* ].
*s-null* := *list* & *s-list*.
```

Notational conventions

 Lists are not available as a built-in data type; abbreviatory notation in tdl:

< a, b > \equiv [FIRST a, REST [FIRST b, REST *null*]]

• Underspecified (variable-length) list:

< a ... > \equiv [FIRST a, REST *list*]

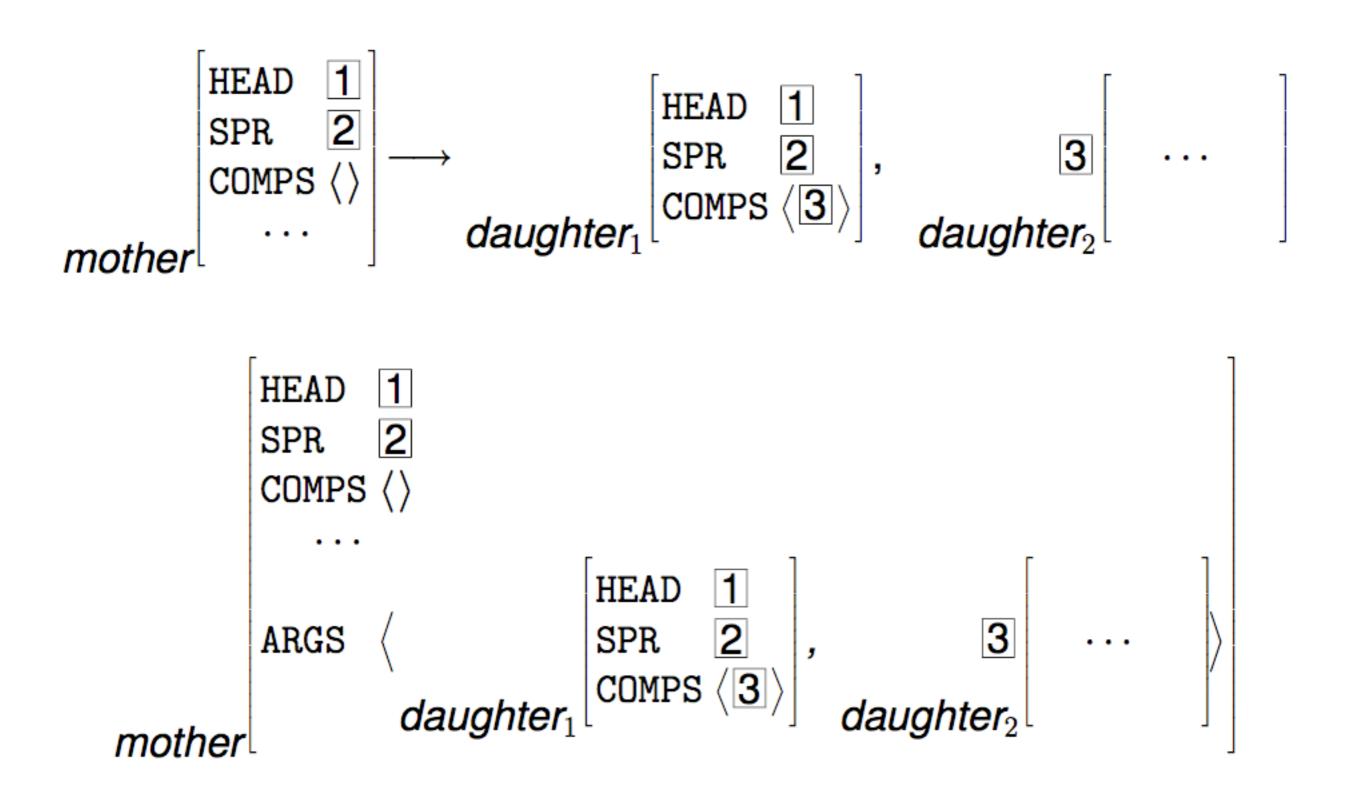
- Difference (open-ended) lists; allow concatenation by unification:
 - <! a !> \equiv [LIST [FIRST a, REST #tail], LAST #tail]

Notational conventions

 strings (e.g., "chased") need no declaration; they are always subtypes of *string*

strings cannot have subtypes, and are (thus) mutually incompatible

Format of grammar rules in the LKB



Overview

- Type hierarchies, inheritance, unification
- Typed feature structures, subsumption, unification
- Type constraints, making typed feature structures wellformed
- Notational conventions
- Grammar rules in the LKB
- Lab 1 questions