# CKY Parsing 

Ling 571
Deep Processing Techniques for NLP January 15, 2014

## Roadmap

- Motivation:
- Parsing (In) efficiency
- Dynamic Programming
- Cocke-Kasami-Younger Parsing Algorithm
- Chomsky Normal Form
- Conversion
- CKY Algorithm
- Parsing by tabulation


## Repeated Work

- Top-down and bottom-up parsing both lead to repeated substructures
- Globally bad parses can construct good subtrees
- But overall parse will fail
- Require reconstruction on other branch
- No static backtracking strategy can avoid
- Efficient parsing techniques require storage of shared substructure
- Typically with dynamic programming
- Example: a flight from Indianapolis to Houston on TWA


## Parsing w/Dynamic Programming

- Avoids repeated work
- Allows implementation of (relatively) efficient parsing algorithms
- Polynomial time in input length
- Typically cubic ( $n^{3}$ ) or less
- Several different implementations
- Cocke-Kasami-Younger (CKY) algorithm
- Earley algorithm
- Chart parsing


## Dynamic Programming in CKY

- Key idea:
- For a parse spanning substring [i,j], there exists some $k$ such there are parses spanning $[i, k]$ and $[k, j]$
- We can construct parses for whole sentence by building up from these stored partial parses


## Dynamic Programming in CKY

- Key idea:
- For a parse spanning substring [i,j], there exists some $k$ such there are parses spanning $[i, k]$ and $[k, j]$
- We can construct parses for whole sentence by building up from these stored partial parses
- So,
- To have a rule A -> B C in [i,j],
- We must have $B$ in $[i, k]$ and $C$ in $[k, j]$, for some $i<k<j$
- CNF grammar forces this for all $j>i+1$


## CKY

- Given an input string $S$ of length $n$,
- Build table $(n+1) \times(n+1)$
- Indexes correspond to inter-word positions
- E.g., 0 Book 1 That 2 Flight 3


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- Build table $(n+1) \times(n+1)$
- Indexes correspond to inter-word positions
- E.g., 0 Book 1 That 2 Flight 3
- Cells [i,j] contain sets of non-terminals of ALL constituents spanning i,j
- $[j-1, j]$ contains pre-terminals
- If $[0, n]$ contains Start, the input is recognized


## CKY Table

- Book the flight through Houston



## CKY Algorithm

function CKY-PARSE(words, grammar) returns table
for $j \leftarrow$ from 1 to LENGTH(words) do table $[j-1, j] \leftarrow\{A \mid A \rightarrow$ words $[j] \in$ grammar $\}$ for $i \leftarrow$ from $j-2$ downto 0 do

$$
\begin{aligned}
& \text { for } k \leftarrow i+1 \text { to } j-1 \text { do } \\
& \qquad \begin{array}{c}
\text { table }[i, j] \leftarrow \text { table }[i, j] \cup \\
\{A \mid A \rightarrow B C \in \text { grammar, }, \\
B \in \text { table }[i, k], \\
C \in \text { table }[k, j]\}
\end{array}
\end{aligned}
$$

## CKY Parsing

- Table fills:
- Column-by-column
- Left-to-right
- Bottom-to-top
- Why?


## CKY Parsing

- Table fills:
- Column-by-column
- Left-to-right
- Bottom-to-top
- Why?
- Necessary info available (below and left)
- Allows online sentence analysis
- Works across input string as it arrives


## Filling a CKY cell



| $\mathscr{L}_{1}$ Grammar | $\mathscr{L}_{1}$ in CNF |
| :---: | :---: |
| $S \rightarrow N P V P$ | $S \rightarrow N P V P$ |
| $S \rightarrow A u x N P V P$ | $S \rightarrow X 1 V P$ |
|  | X1 $\rightarrow$ Aux NP |
| $S \rightarrow V P$ | $S \rightarrow$ book \| include | prefer |
|  | $S \rightarrow \operatorname{Verb}$ NP |
|  | $S \rightarrow X 2 P P$ |
|  | $S \rightarrow V \mathrm{Verb} P \mathrm{P}$ |
|  | $S \rightarrow V P P P$ |
| $N P \rightarrow$ Pronoun | $N P \rightarrow I \mid$ she $\mid$ me |
| $N P \rightarrow$ Proper-Noun | $N P \rightarrow$ TWA $\mid$ Houston |
| $N P \rightarrow$ Det Nominal | $N P \rightarrow$ Det Nominal |
| Nominal $\rightarrow$ Noun | Nominal $\rightarrow$ book $\mid$ flight $\mid$ meal $\mid$ money |
| Nominal $\rightarrow$ Nominal Noun | Nominal $\rightarrow$ Nominal Noun |
| Nominal $\rightarrow$ Nominal PP | Nominal $\rightarrow$ Nominal PP |
| $V P \rightarrow$ Verb | $V P \rightarrow$ book \| include | prefer |
| $V P \rightarrow$ Verb $N P$ | $V P \rightarrow$ Verb $N P$ |
| $V P \rightarrow \operatorname{Verb} N P P P$ | $V P \rightarrow X 2 P P$ |
|  | $X 2 \rightarrow \operatorname{Verb} N P$ |
| $V P \rightarrow$ Verb $P P$ | $V P \rightarrow V \operatorname{Verb} P P$ |
| $V P \rightarrow V P P P$ | $V P \rightarrow V P P P$ |
| $P P \rightarrow$ Preposition NP | $P P \rightarrow$ Preposition $N P$ |

0 Book 1 the 2 flight 3 through 4 Houston 5

| Book | the | Flight | through | Houston |
| :--- | :--- | :--- | :--- | :--- |
| NN, VB, <br> Nominal, VP, S <br> $[0,1]$ |  |  |  |  |
|  |  |  |  |  |
|  |  |  |  |  |
|  |  |  |  |  |

## 0 Book 1 the 2 flight 3 through 4 Houston 5

| Book | the | Flight | through | Houston |
| :--- | :--- | :--- | :--- | :--- |
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|  |  |  |  |  |



## 0 Book 1 the 2 flight 3 throught 4 Houston 5

| Book | the | Flight | Through | Houston |
| :--- | :--- | :--- | :--- | :--- |
| NN, VB, <br> Nominal, VP, S <br> $[0,1]$ | [0,2] |  |  |  |
|  | Det <br> $[1,2]$ |  |  |  |
|  |  |  |  |  |
|  |  |  |  |  |

0 Book 1 the 2 flight 3 through 4 Houston 5
$\left.\begin{array}{|l|l|l|l|l|}\hline \text { Book } & \text { the } & \text { Flight } & \text { through } & \text { Houston } \\ \hline \begin{array}{l}\text { NN, VB, } \\ \text { Nominal, VP, S } \\ {[0,1]}\end{array} & & & & \\ \hline & {[0,2]}\end{array}\right)$

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## 0 Book 1 the 2 flight 3 through 4 Houston 5

| Book | the | Flight | Through | Houston |
| :--- | :--- | :--- | :--- | :--- |
| NN, VB, <br> Nominal, VP, S <br> $[0,1]$ |  | S, VP, X2 |  |  |
|  | $[0,2]$ | $[0,3]$ |  |  |
|  | Det <br> $[1,2]$ | NP <br> $[1,3]$ |  |  |
|  |  |  | NN, Nominal <br> $[2,3]$ |  |

## 0 Book 1 the 2 flight 3 through 4 Houston 5

| Book | the | Flight | Through | Houston |
| :---: | :---: | :---: | :---: | :---: |
| NN, VB, |  | S, VP, X2 |  |  |
|  | [0,2] | [0,3] |  |  |
|  | $\begin{aligned} & \text { Det } \\ & {[1,2]} \end{aligned}$ | $\begin{aligned} & \text { NP } \\ & {[1,3]} \end{aligned}$ |  |  |
|  |  | NN, Nominal [2,3] |  |  |
|  |  |  | $\begin{aligned} & \text { Prep } \\ & {[3,4]} \end{aligned}$ |  |
|  |  |  |  |  |

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| Book | the | Flight | Through | Houston |
| :--- | :--- | :--- | :--- | :--- |
| NN, VB, <br> Nominal, VP, S <br> $[0,1]$ |  | S, VP, X2 |  |  |
|  | $[0,2]$ | $[0,3]$ | $[0,4]$ |  |
|  | Det | NP |  |  |
| $[1,2]$ |  | NN, Nominal <br> $[2,3]$ | $[2,4]$ |  |

## 0 Book 1 the 2 flight 3 through 4 Houston 5

| Book | the | Flight | Through | Houston |
| :---: | :---: | :---: | :---: | :---: |
| NN, VB, Nominal, VP, S $[0,1]$ |  | S, VP, X2 |  |  |
|  | $[0,2]$ |  | [0,4] |  |
|  | $\begin{aligned} & \text { Det } \\ & {[1,2]} \end{aligned}$ | $\begin{aligned} & N P \\ & {[1,3]} \end{aligned}$ | [1,4] |  |
|  |  | NN, Nominal [2,3] | [2,4] |  |
|  |  |  | $\begin{aligned} & \text { Prep } \\ & {[3,4]} \end{aligned}$ |  |
|  |  |  |  | NNP, NP <br> [4,5] |

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| :---: | :---: | :---: | :---: | :---: |
| NN, VB, Nominal, VP, S $[0,1]$ | [0,2] | $\begin{aligned} & \mathrm{S}, \mathrm{VP}, \mathrm{X} 2 \\ & {[0,3]} \end{aligned}$ | [0,4] |  |
|  | $\begin{aligned} & \text { Det } \\ & {[1,2]} \end{aligned}$ | $\begin{aligned} & N P \\ & {[1,3]} \end{aligned}$ | [1,4] |  |
|  |  | NN, Nominal $[2,3]$ | [2,4] |  |
|  |  |  | $\begin{aligned} & \text { Prep } \\ & {[3,4]} \end{aligned}$ | $\begin{aligned} & \text { PP } \\ & {[3,5]} \end{aligned}$ |
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| :--- | :--- | :--- | :--- | :--- |
| NN, VB, <br> Nominal, VP, S <br> $[0,1]$ |  | S, VP, X2 |  |  |
|  | $[0,2]$ | $[0,3]$ | $[0,4]$ |  |
|  | Det |  |  |  |
| $[1,2]$ | NP | $[1,3]$ | $[1,4]$ |  |
|  |  |  |  |  |

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| NN, VB, <br> Nominal, VP, S <br> $[0,1]$ |  | S, VP, X2 |  |  |
|  | $[0,2]$ | $[0,3]$ | $[0,4]$ |  |
|  | Det | NP |  | NP |
|  | $[1,2]$ | $[1,3]$ | $[1,4]$ | $[1,5]$ |
|  |  |  |  | NN, Nominal |
| $[2,3]$ | $[2,4]$ | Nominal <br>  <br>  |  |  |
|  |  |  | Prep | PP |

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| Book | the | Flight | Through | Houston |
| :---: | :---: | :---: | :---: | :---: |
| NN, VB, Nominal, VP, S [0,1] | [0,2] | $\begin{aligned} & \mathrm{S}, \mathrm{VP}, \mathrm{X} 2 \\ & {[0,3]} \end{aligned}$ | [0,4] | $\begin{aligned} & \mathrm{S}, \mathrm{VP}, \mathrm{X} 2 \\ & {[0,5]} \end{aligned}$ |
|  | $\begin{aligned} & \text { Det } \\ & {[1,2]} \end{aligned}$ | $\begin{aligned} & N P \\ & {[1,3]} \end{aligned}$ | [1,4] | $\begin{aligned} & N P \\ & {[1,5]} \end{aligned}$ |
|  |  | NN, Nominal $[2,3]$ | [2,4] | Nominal $[2,5]$ |
|  |  |  | $\begin{aligned} & \text { Prep } \\ & {[3,4]} \end{aligned}$ | $\begin{aligned} & \text { PP } \\ & {[3,5]} \end{aligned}$ |
|  |  |  |  | NNP, NP $[4,5]$ |

## Is this a parser?

## From Recognition to Parsing

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- Parsing solution:
- All repeated versions of non-terminals
- Pair each non-terminal with pointers to cells
- Backpointers
- Last step: construct trees from back-pointers in [0, n]


## Filling column 5

| Book | the | flight | through | Houston |
| :---: | :---: | :---: | :---: | :---: |
| S, VP, Verb, Nominal, Noun [0,1] | [0,2] | $\begin{aligned} & S, V P, \times 2 \\ & {[0,3]} \end{aligned}$ | [0,4] | [0,5] |
|  | Det $[1,2]$ | $\begin{array}{\|c} \text { NP } \\ {[1,3]} \\ \hline \end{array}$ | $[1,4]$ | [1,5] |
|  |  | Nominal, Noun $[2,3]$ | [2,4] | Nominal $[2,5]$ |
|  |  |  | Prep $[3,4]$ | [3,5] |
|  |  |  |  | NP, ProperNoun [4,5] |




| Book | the | flight | through | Houston |
| :---: | :---: | :---: | :---: | :---: |
| S, VP, Verb Nominal, Noun $[0,1]$ | [0,2] | $\begin{aligned} & \mathrm{S}, \mathrm{VP}, \mathrm{X2} \\ & {[0,3]} \end{aligned}$ | [0,4] | [0,5] |
|  | Det $[1,2]$ | $[1,3]$ | $[1,4]$ |  |
|  |  | Nominal, Noun $[2,3]$ | [2,4] | Nominal $[2,5]$ |
|  |  |  | Prep $[3,4]$ | $\begin{aligned} & \mathrm{PP} \\ & {[3,5]} \end{aligned}$ |
|  |  |  |  | NP, ProperNoun $[4,5]$ |


| Book | the | flight | through | Houston |
| :---: | :---: | :---: | :---: | :---: |
| S, VP, Verb Nominal, Noun $[0,1]$ | $[0,2]$ | $\begin{aligned} & \hline \mathrm{S}, \\ & \mathrm{VP}, \leftarrow \\ & \times 2 \\ & {[0,3]} \end{aligned}$ | [0,4] | $\begin{array}{r} S_{1}, V P, \times 2 \\ \qquad-S_{2}, \text { VP } \\ \hline \downarrow \\ \hline \end{array}$ |
|  | Det $[1,2]$ | NP $[1,3]$ | $[1,4]$ | $\begin{gathered} \mathrm{NP} \\ {[1,5]} \end{gathered}$ |
|  |  | Nominal, Noun $[2,3]$ | [2,4] |  |
|  |  |  | Prep $[3,4]$ | $[3,5]$ |
|  |  |  |  | NP, Proper- |

## CKY Discussion

- Running time:

$$
O\left(n^{3}\right)
$$

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- Expressiveness:
- As implemented, requires CNF
- Weakly equivalent to original grammar
- Doesn't capture full original structure
- Back-conversion?
- Can do binarization, terminal conversion
- Unit non-terminals require change in CKY


## Parsing Efficiently

- With arbitrary grammars
- Earley algorithm
- Top-down search
- Dynamic programming
- Tabulated partial solutions
- Some bottom-up constraints


## Earley Parsing

- Avoid repeated work/recursion problem
- Dynamic programming
- Store partial parses in "chart"
- Compactly encodes ambiguity
- $O\left(N^{3}\right)$


## Earley Parsing

- Avoid repeated work/recursion problem
- Dynamic programming
- Store partial parses in "chart"
- Compactly encodes ambiguity - $O\left(N^{3}\right)$
- Chart entries:
- Subtree for a single grammar rule
- Progress in completing subtree
- Position of subtree wrt input


## Earley Algorithm

- First, left-to-right pass fills out a chart with N+1 states
- Think of chart entries as sitting between words in the input string, keeping track of states of the parse at these positions
- For each word position, chart contains set of states representing all partial parse trees generated to date. E.g. chart[0] contains all partial parse trees generated at the beginning of the sentence


## Chart Entries

## Represent three types of constituents:

- predicted constituents
- in-progress constituents
- completed constituents


## Parse Progress

- Represented by Dotted Rules
- Position of •indicates type of constituent
- ${ }_{0}$ Book $_{1}$ that ${ }_{2}$ flight $_{3}$
- $S \rightarrow$ •VP, [0,0] (predicted)


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- S $\rightarrow$ VP, [0,0] (predicted)
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- VP $\rightarrow$ V NP •, [0,3] (completed)
- $[x, y]$ tells us what portion of the input is spanned so far by this rule


## Parse Progress

- Represented by Dotted Rules
- Position of •indicates type of constituent
- oBook ${ }_{1}$ that ${ }_{2}$ flight ${ }_{3}$
- $S \rightarrow \cdot V P,[0,0]$ (predicted)
- NP $\rightarrow$ Det • Nom, $[1,2]$ (in progress)
- VP $\rightarrow$ V NP •, [0,3] (completed)
- $[\mathrm{x}, \mathrm{y}]$ tells us what portion of the input is spanned so far by this rule
- Each State $\mathrm{s}_{\mathrm{i}}$ : <dotted rule>, [<back pointer>,<current position>]


## ${ }_{0}$ Book $_{1}$ that ${ }_{2}$ flight 3

## $S \rightarrow \cdot V P,[0,0]$

- First 0 means $S$ constituent begins at the start of input
- Second 0 means the dot here, too
- So, this is a top-down prediction


## ${ }_{0}$ Book $_{1}$ that ${ }_{2}$ flight 3

$S \rightarrow \cdot V P,[0,0]$

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- Second 0 means the dot here too
- So, this is a top-down prediction

NP $\rightarrow$ Det •Nom, $[1,2]$

- the NP begins at position 1
- the dot is at position 2
- so, Det has been successfully parsed
- Nom predicted next


# ${ }_{0}$ Book $_{1}$ that ${ }_{2}$ flight 3 (continued) 

$$
\mathrm{VP} \rightarrow \mathrm{~V} \mathrm{NP} \mathrm{\cdot,} \mathrm{[0,3]}
$$

- Successful VP parse of entire input



## Successful Parse

- Final answer found by looking at last entry in chart


## Successful Parse

- Final answer found by looking at last entry in chart
- If entry resembles $S \rightarrow \alpha \cdot[0, N]$ then input parsed successfully
- Chart will also contain record of all possible parses of input string, given the grammar


## Parsing Procedure for the Earley Algorithm

- Move through each set of states in order, applying one of three operators to each state:
- predictor: add predictions to the chart
- scanner: read input and add corresponding state to chart
- completer: move dot to right when new constituent found


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- Move through each set of states in order, applying one of three operators to each state:
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- Results (new states) added to current or next set of states in chart

No backtracking and no states removed: keep complete history of parse

## States and State Sets

- Dotted Rule si represented as <dotted rule>, [<back pointer>, <current position>]
- State Set $\mathbf{S}_{\mathrm{j}}$ to be a collection of states $\mathrm{s}_{\mathrm{i}}$ with the same <current position>.


## Earley Algorithm from Book

function EARLEY-PARSE(words, grammar) returns chart

```
ENQUEUE(( }\gamma->\bulletS,[0,0]),\operatorname{chart[0])
for }i\leftarrow\mathrm{ from 0 to LENGTH(words) do
    for each state in chart[i] do
        if INCOMPLETE?(state) and
            NEXT-CAT(state) is not a part of speech then
        PREDICTOR(state)
        elseif INCOMPLETE?(state) and
            NEXT-CAT(state) is a part of speech then
        SCANNER(state)
        else
        COMPLETER(state)
    end
end
return(chart)
```


## Earley Algorithm from Book

procedure PREDICTOR $((A \rightarrow \alpha \bullet B \beta,[i, j]))$
for each $(B \rightarrow \gamma)$ in GRAMMAR-RULES-FOR( $B$, grammar) do EnQueue $((B \rightarrow \bullet \gamma,[j, j]), \operatorname{chart}[j])$ end
procedure $\operatorname{SCANNER}((A \rightarrow \alpha \bullet B \beta,[i, j]))$
if $\mathrm{B} \subset$ PARTS-OF-SPEECH(word[j]) then
$\operatorname{ENQUEUE}((B \rightarrow \operatorname{word}[j],[j, j+1]), \operatorname{chart}[j+1])$
procedure COMPLETER $((B \rightarrow \quad \gamma \quad$ • $[j, k]))$
for each $(A \rightarrow \alpha \bullet B \beta,[i, j])$ in $\operatorname{chart}[j]$ do
EnQueue $((A \rightarrow \alpha B \bullet \beta,[i, k])$, chart $[k])$
end

## Earley Algorithm (simpler!)

1. Add Start $\rightarrow$ S, $[0,0]$ to state set 0

Let $\mathrm{i}=1$
2. Predict all states you can, adding new predictions to state set 0
3. Scan input word $i$-add all matched states to state set $\mathrm{S}_{\mathrm{i}}$. Add all new states produced by Complete to state set $S_{i}$ Add all new states produced by Predict to state set $\mathrm{S}_{\mathrm{i}}$ Let $\mathrm{i}=\mathrm{i}+1$
Unless $\mathrm{i}=n$, repeat step 3 .
4. At the end, see if state set $n$ contains Start $\rightarrow \mathrm{S} \cdot,[0, \mathrm{n}]$

## 3 Main Sub-Routines of Earley Algorithm

- Predictor: Adds predictions into the chart.
- Completer: Moves the dot to the right when new constituents are found.
- Scanner: Reads the input words and enters states representing those words into the chart.


## Predictor

- Intuition: create new state for top-down prediction of new phrase.


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- Intuition: create new state for top-down prediction of new phrase.
- Applied when non part-of-speech nonterminals are to the right of a dot: $\mathbf{S} \rightarrow \bullet$ VP [0,0]
- Adds new states to current chart
- One new state for each expansion of the nonterminal in the grammar
VP $\rightarrow$ • V [0,0]
VP $\rightarrow$ • V NP [0,0]


## Predictor

- Intuition: create new state for top-down prediction of new phrase.
- Applied when non part-of-speech nonterminals are to the right of a dot: $\mathbf{S} \rightarrow \bullet$ VP [0,0]
- Adds new states to current chart
- One new state for each expansion of the nonterminal in the grammar VP $\rightarrow$ • V [0,0] VP $\rightarrow$ • V NP [0,0]
Formally:



## Chart[0]

S0

[0,0]
Dummy start state

Note that given a grammar, these entries are the same for all inputs; they can be pre-loaded.

## Chart[0]

| S0 | $\gamma \rightarrow \bullet S$ | $[0,0]$ | Dummy start state |
| :--- | :--- | ---: | ---: |
| S1 | $S \rightarrow \bullet N P V P$ | $[0,0]$ | Predictor |
| S2 | $S \rightarrow \bullet A u x N P V P$ | $[0,0]$ | Predictor |
| S3 | $S \rightarrow \bullet V P$ | $[0,0]$ | Predictor |
|  |  |  |  |
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|  |  |  |  |
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Note that given a grammar, these entries are the same for all inputs; they can be pre-loaded.

## Chart[0]

| S0 | $\gamma \rightarrow \bullet S$ | $[0,0]$ | Dummy start state |
| :--- | :--- | :--- | ---: |
| S1 | $S \rightarrow \bullet N P V P$ | $[0,0]$ | Predictor |
| S2 | $S \rightarrow \bullet$ Aux NPVP | $[0,0]$ | Predictor |
| S3 | $S \rightarrow \bullet V P$ | $[0,0]$ | Predictor |
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| S5 | $N P \rightarrow$ Proper -Noun | $[0,0]$ | Predictor |
| S6 | $N P \rightarrow \bullet$ Det Nominal | $[0,0]$ | Predictor |
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| S6 | $N P \rightarrow \bullet$ Det Nominal | $[0,0]$ | Predictor |
| S7 | $V P \rightarrow \bullet$ Verb | $[0,0]$ | Predictor |
| S8 | $V P \rightarrow \bullet$ Verb $N P$ | $[0,0]$ | Predictor |
| S9 | $V P \rightarrow \bullet$ Verb $N P P P$ | $[0,0]$ | Predictor |
| S10 | $V P \rightarrow \bullet$ Verb $P P$ | $[0,0]$ | Predictor |
| S11 | $V P \rightarrow \bullet V P P P$ | $[0,0]$ | Predictor |

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- Formally:

$$
\begin{aligned}
& S_{j}: A \rightarrow \alpha \cdot B \beta,[i, j] \\
& S_{j+1}: A \rightarrow \alpha B \cdot \beta,[i, j+1]
\end{aligned}
$$

## Completer

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- VP $\rightarrow$ V NP • 0,3$]$
- Formally: $\mathrm{S}_{\mathrm{K}}: \mathrm{B} \rightarrow \boldsymbol{\delta}^{\delta} \cdot[j, \mathrm{k}]$
$S_{\text {whe }}: A \rightarrow \alpha \rightarrow S_{j}: A \rightarrow \alpha,[i, k \in, \beta,[i, j]$.


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## Chart[1]

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S12 Verb $\rightarrow$ book •
[0,1]
Scanner

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| S12 | Ver $b \rightarrow$ book $\bullet$ | $[0,1]$ | Scanner |
| :--- | :--- | :--- | ---: |
| S13 | $V P \rightarrow$ Verb $\bullet$ | $[0,1]$ | Completer |
| S14 | $V P \rightarrow$ Verb $\bullet N P$ | $[0,1]$ | Completer |
| S15 | $V P \rightarrow$ Ver $\bullet N P P P$ | $[0,1]$ | Completer |
| S16 | $V P \rightarrow$ Ver $\bullet P P$ | $[0,1]$ | Completer |

## Chart[1]

| S12 | Ver $b \rightarrow$ book $\bullet$ | $[0,1]$ | Scanner |
| :--- | :--- | :--- | ---: |
| S13 | $V P \rightarrow$ Ver $\bullet$ | $[0,1]$ | Completer |
| S14 | $V P \rightarrow V e r b \bullet N P$ | $[0,1]$ | Completer |
| S15 | $V P \rightarrow \operatorname{Ver} \bullet \bullet N P P P$ | $[0,1]$ | Completer |
| S16 | $V P \rightarrow \operatorname{Ver} \bullet \bullet P P$ | $[0,1]$ | Completer |
| S17 | $S \rightarrow V P \bullet$ | $[0,1]$ | Completer |
| S18 | $V P \rightarrow V P \bullet P P$ | $[0,1]$ | Completer |

## Chart[1]

| S12 | Verb $\rightarrow$ book $\bullet$ | $[0,1]$ | Scanner |
| :--- | :--- | :--- | ---: |
| S13 | $V P \rightarrow$ Verb $\bullet$ | $[0,1]$ | Completer |
| S14 | $V P \rightarrow$ Verb $\bullet N P$ | $[0,1]$ | Completer |
| S15 | $V P \rightarrow$ Verb $\bullet N P P P$ | $[0,1]$ | Completer |
| S16 | $V P \rightarrow$ Verb $\bullet P P$ | $[0,1]$ | Completer |
| S17 | $S \rightarrow V P \bullet$ | $[0,1]$ | Completer |
| S18 | $V P \rightarrow V P \bullet P P$ | $[0,1]$ | Completer |
| S19 | $N P \rightarrow \bullet$ Pronoun | $[1,1]$ | Predictor |
| S20 | $N P \rightarrow \bullet$ Proper-Noun | $[1,1]$ | Predictor |
| S21 | $N P \rightarrow \bullet$ Det Nominal | $[1,1]$ | Predictor |
| S22 | $P P \rightarrow \bullet$ Prep NP | $[1,1]$ | Predictor |

## Prediction of Next Rule

- When $\mathrm{VP} \rightarrow \mathrm{V}$ • is itself processed by the Completer, $\mathrm{S} \rightarrow \mathrm{VP}$ • is added to Chart[1] since VP is a left corner of $S$
- Last few rules in Chart[1] are added by Predictor when VP $\rightarrow \mathrm{V} \bullet \mathrm{NP}$ is processed
- And so on....


## Charts[2] and [3]

## Charts[2] and [3]

S23 Det $\rightarrow$ that $\bullet$
[1,2]

## Charts[2] and [3]

| S23 | Det $\rightarrow$ that $\bullet$ | $[1,2]$ |
| :--- | :--- | :--- |
| S24 | $N P \rightarrow$ Det $\bullet$ Nominal | $[1,2]$ |

Scanner
Completer

## Charts[2] and [3]

| S23 | Det $\rightarrow$ that $\bullet$ | $[1,2]$ | Scanner |
| :--- | :--- | ---: | ---: |
| S24 | NP $\rightarrow$ Det $\bullet$ Nominal | $[1,2]$ | Completer |
| S25 | Nominal $\rightarrow$ Noun | $[2,2]$ | Predictor |
| S26 | Nominal $\rightarrow \bullet$ Nominal Noun | $[2,2]$ | Predictor |
| S27 | Nominal $\rightarrow \bullet$ Nominal PP | $[2,2]$ | Predictor |

## Charts[2] and [3]

| S23 | Det $\rightarrow$ that $\bullet$ | $[1,2]$ | Scanner |
| :--- | :--- | ---: | ---: |
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| S28 | Noun $\rightarrow$ flight $\bullet$ | $[2,3]$ | Scanner |

## Charts[2] and [3]

| S23 | Det $\rightarrow$ that $\bullet$ | $[1,2]$ | Scanner |
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| S26 | Nominal $\rightarrow \bullet$ Nominal Noun | $[2,2]$ | Predictor |
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| S28 | Noun $\rightarrow$ flight $\bullet$ | $[2,3]$ | Scanner |
| S29 | Nominal $\rightarrow$ Noun $\bullet$ | $[2,3]$ | Completer |
| S30 | NP $\rightarrow$ Det Nominal $\bullet$ | $[1,3]$ | Completer |
| S31 | Nominal $\rightarrow$ Nominal $\bullet$ Noun | $[2,3]$ | Completer |
| S32 | Nominal $\rightarrow$ Nominal $\bullet P P$ | $[2,3]$ | Completer |
| S33 | $V P \rightarrow$ Verb NP $\bullet$ | $[0,3]$ | Completer |
| S34 | $V P \rightarrow$ Verb NP $\bullet P P$ | $[0,3]$ | Completer |
| S35 | $P P \rightarrow \bullet$ Prep NP | $[3,3]$ | Predictor |
| S36 | $S \rightarrow$ VP $\bullet$ | $[0,3]$ | Completer |

## How do we retrieve the parses at the end?

- Augment the Completer to add pointers to prior states it advances as a field in the current state
- i.e. what state did we advance here?
- Read the pointers back from the final state
- What about ambiguity?
- What about ambiguity?
- CKY/Earley can represent it
- What about ambiguity?
- CKY/Earley can represent it
- Can't resolve it

