# **CKY** Parsing

Ling571 Deep Processing Approaches to NLP January 11, 2017

#### Roadmap

- Motivation:
  - Inefficiencies of parsing-as-search
- Strategy: Dynamic Programming
- Chomsky Normal Form
  - Weak and strong equivalence
- CKY parsing algorithm

# **Bottom-Up Parsing**

- Try to find all trees that span the input
  - Start with input string
    - Book that flight.
  - Use all productions with current subtree(s) on RHS
    - E.g., N → Book; V → Book
  - Stop when spanned by S (or no more rules apply)

# Book that flight

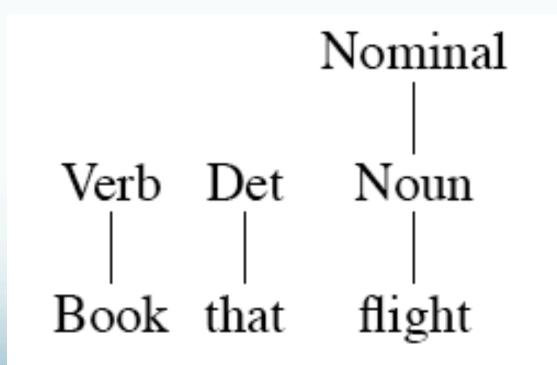
Speech and Language Processing -

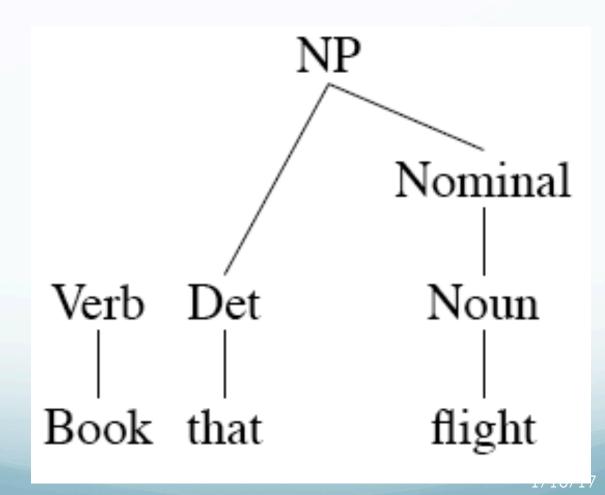
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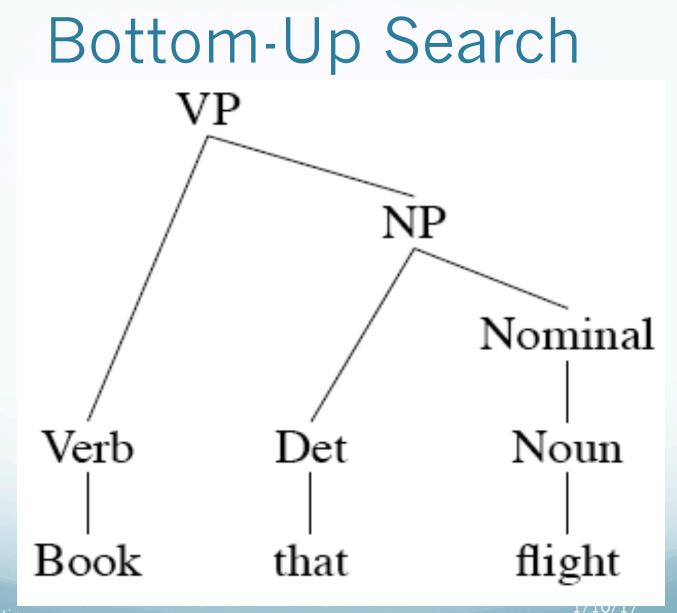
# Verb Det Noun Book that flight

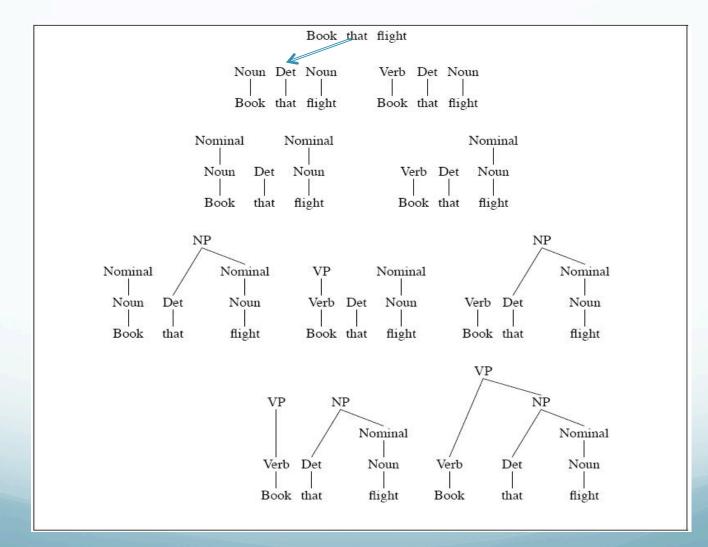
Speech and Language Processing -

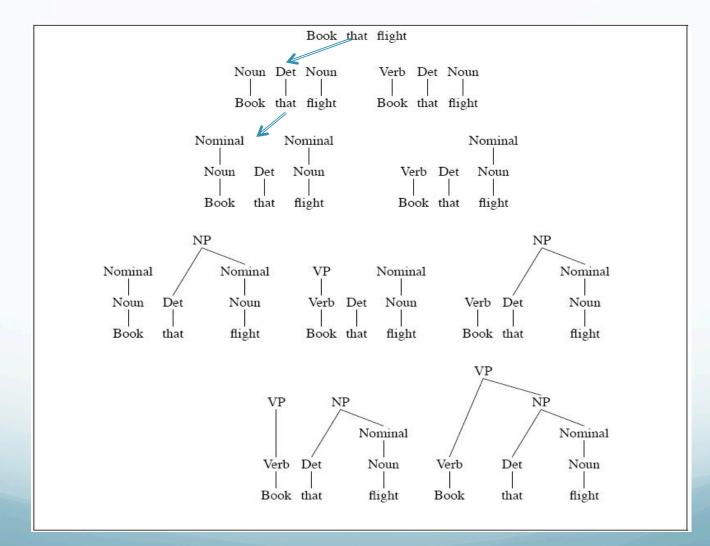
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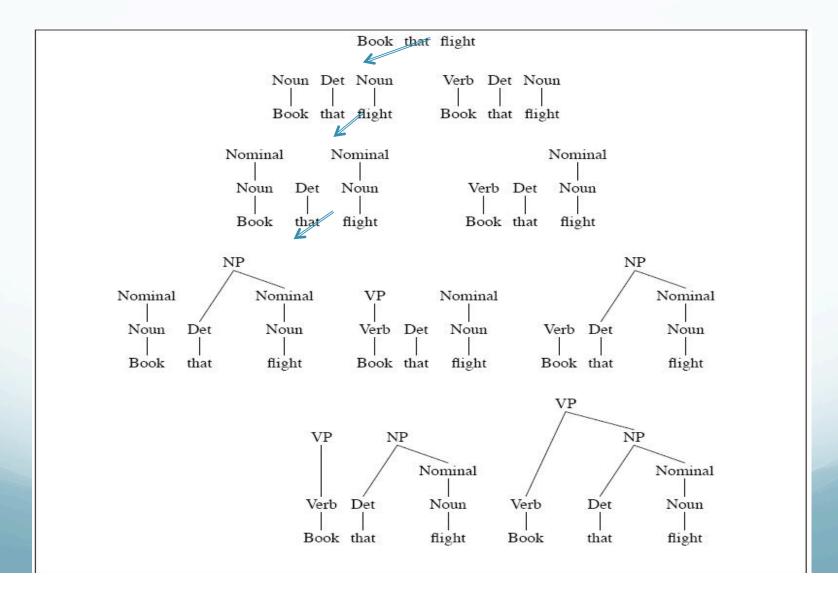












# Pros and Cons of Bottom-Up Search

#### • Pros:

- Will not explore trees that don't match input
- Recursive rules less problematic
- Useful for incremental/ fragment parsing
- Cons:
  - Explore subtrees that will not fit full sentences

## Parsing Challenges

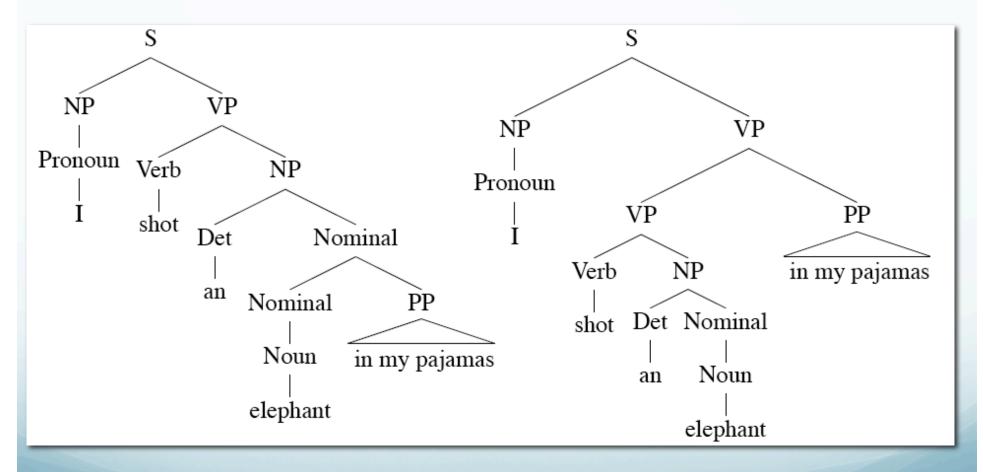
- Ambiguity
- Repeated substructure

• Recursion

# Parsing Ambiguity

- Many sources of parse ambiguity
  - Lexical ambiguity
    - Book/N; Book/V
  - Structural ambiguity: Main types:
    - Attachment ambiguity
      - Constituent can attach in multiple places
        - I shot an elephant in my pyjamas.
    - Coordination ambiguity
      - Different constituents can be conjoined
        - Old men and women

# Ambiguity



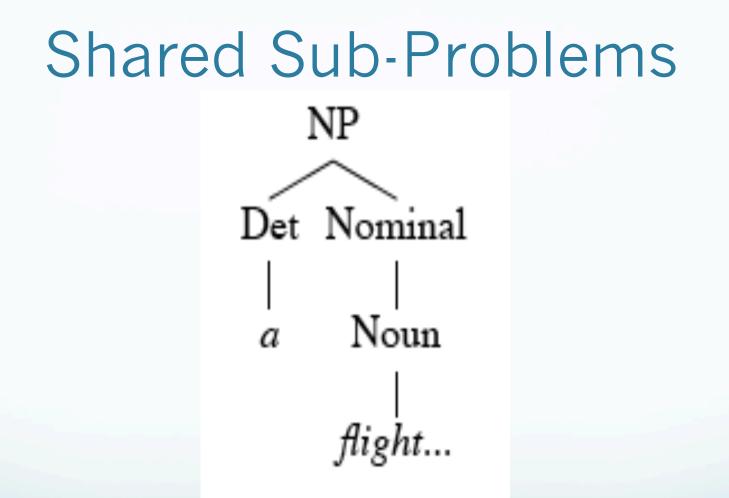
Speech and Language Processing -

# Disambiguation

- Global ambiguity:
  - Multiple complete alternative parses
  - Need strategy to select correct one
    - Approaches exploit other information
      - Statistical
        - Some prepositional structs more likely to attach high/low
        - Some phrases more likely, e.g., (old (men and women))
      - Semantic
      - Pragmatic
        - E.g., elephants and pyjamas
  - Alternatively, keep all
- Local ambiguity:
  - Ambiguity in subtree, resolved globally

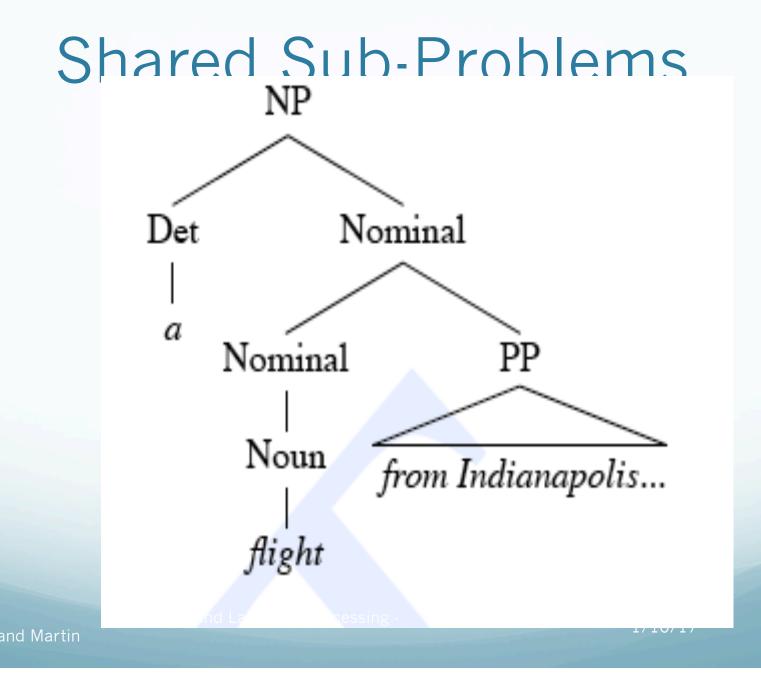
#### Repeated Work

- Top-down and bottom-up parsing both lead to repeated substructures
  - Globally bad parses can construct good subtrees
    - But overall parse will fail
    - Require reconstruction on other branch
  - No static backtracking strategy can avoid
- Efficient parsing techniques require storage of shared substructure
  - Typically with dynamic programming
- Example: a flight from Indianapolis to Houston on TWA

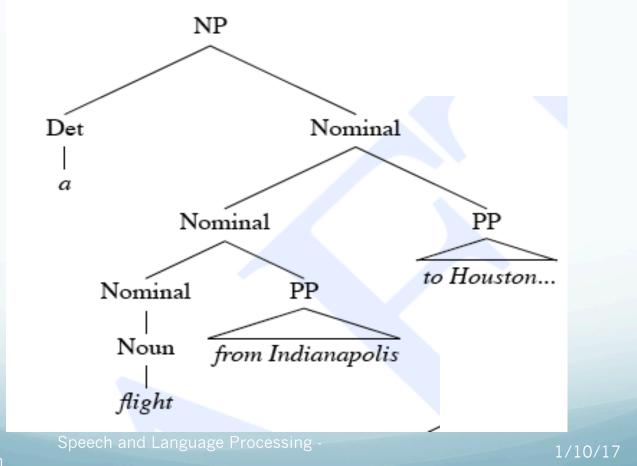


Speech and Language Processing -

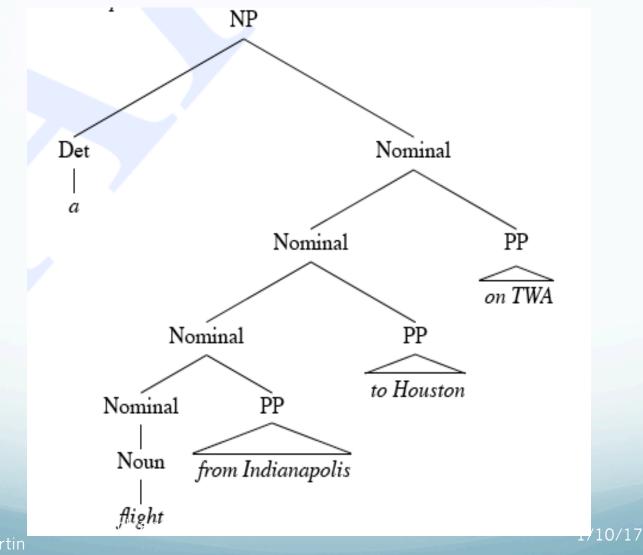
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#### Shared Sub-Problems



#### Shared Sub-Problems



#### Recursion

- Many grammars have recursive rules
  - E.g., S → S Conj S
- In search approaches, recursion is problematic
  - Can yield infinite searches
    - Esp., top-down

# Dynamic Programming

- Challenge: Repeated substructure → Repeated work
- Insight:
  - Global parse composed of parse substructures
  - Can record parses of substructures
- Dynamic programming avoids repeated work by tabulating solutions to subproblems
  - Here, stores subtrees

# Parsing w/Dynamic Programming

- Avoids repeated work
- Allows implementation of (relatively) efficient parsing algorithms
  - Polynomial time in input length
    - Typically cubic (  $n^3$  ) or less
- Several different implementations
  - Cocke-Kasami-Younger (CKY) algorithm
  - Earley algorithm
  - Chart parsing

# Chomsky Normal Form (CNF)

- CKY parsing requires grammars in CNF
- Chomsky Normal Form
  - All productions of the form:
    - $A \rightarrow B C$ , or
    - A → a
- However, most of our grammars are not of this form
  - E.g., S → Wh-NP Aux NP VP
- Need a general conversion procedure
  - Any arbitrary grammar can be converted to CNF

#### Grammar Equivalence and Form

• Grammar equivalence

- Weak: Accept the same language, May produce different analyses
- Strong: Accept same language, Produce same structure

## **CNF** Conversion

- Three main conditions:
  - Hybrid rules:
    - INF-VP → to VP
  - Unit productions:

•  $A \rightarrow B$ 

- Long productions:
  - $A \rightarrow B C D$

## **CNF** Conversion

- Hybrid rule conversion:
  - Replace all terminals with dummy non-terminals
  - E.g., INF-VP → to VP
    - INF-VP  $\rightarrow$  TO VP; TO  $\rightarrow$  to
- Unit productions:
  - Rewrite RHS with RHS of all derivable non-unit productions
    - If  $A \Rightarrow B$  and  $B \rightarrow w$ , then add  $A \rightarrow w$

## **CNF** Conversion

- Long productions:
  - Introduce new non-terminals and spread over rules
  - S  $\rightarrow$  Aux NP VP
    - $S \rightarrow X1 VP; X1 \rightarrow Aux NP$
- For all non-conforming rules,
  - Convert terminals to dummy non-terminals
  - Convert unit productions
  - Binarize all resulting rules

$\mathscr{L}_1$ Grammar	$\mathscr{L}_1$ in CNF
$S \rightarrow NP VP$	$S \rightarrow NP VP$
$S \rightarrow Aux NP VP$	$S \rightarrow X1 VP$
	$X1 \rightarrow Aux NP$
$S \rightarrow VP$	$S \rightarrow book \mid include \mid prefer$
	$S \rightarrow Verb NP$
	$S \rightarrow X2 PP$
	$S \rightarrow Verb PP$
	$S \rightarrow VP PP$
$NP \rightarrow Pronoun$	$NP \rightarrow I \mid she \mid me$
$NP \rightarrow Proper-Noun$	$NP \rightarrow TWA \mid Houston$
$NP \rightarrow Det Nominal$	$NP \rightarrow Det Nominal$
$Nominal \rightarrow Noun$	Nominal $\rightarrow$ book   flight   meal   money
$Nominal \rightarrow Nominal Noun$	Nominal $\rightarrow$ Nominal Noun
Nominal $\rightarrow$ Nominal PP	Nominal $\rightarrow$ Nominal PP
$VP \rightarrow Verb$	$VP \rightarrow book \mid include \mid prefer$
$VP \rightarrow Verb NP$	$VP \rightarrow Verb NP$
$VP \rightarrow Verb NP PP$	$VP \rightarrow X2 PP$
	$X2 \rightarrow Verb NP$
$VP \rightarrow Verb PP$	$VP \rightarrow Verb PP$
$VP \rightarrow VP PP$	$VP \rightarrow VP PP$
$PP \rightarrow Preposition NP$	$PP \rightarrow Preposition NP$

# **CKY** Parsing

- Cocke-Kasami-Younger parsing algorithm:
  - (Relatively) efficient bottom-up parsing algorithm based on tabulating substring parses to avoid repeated work
  - Approach:
    - Use a CNF grammar
    - Build an (n+1) x (n+1) matrix to store subtrees
      - Upper triangular portion
    - Incrementally build parse spanning whole input string

# Dynamic Programming in CKY

- Key idea:
  - For a parse spanning substring [i,j], there exists some k such there are parses spanning [i,k] and [k,j]
    - We can construct parses for whole sentence by building up from these stored partial parses
- So,
  - To have a rule  $A \rightarrow B C$  in [i,j],
    - We must have B in [i,k] and C in [k,j], for some i<k<j</li>
      - CNF grammar forces this for all j>i+1